

# Introduction to Artificial Intelligence



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**COMP307**

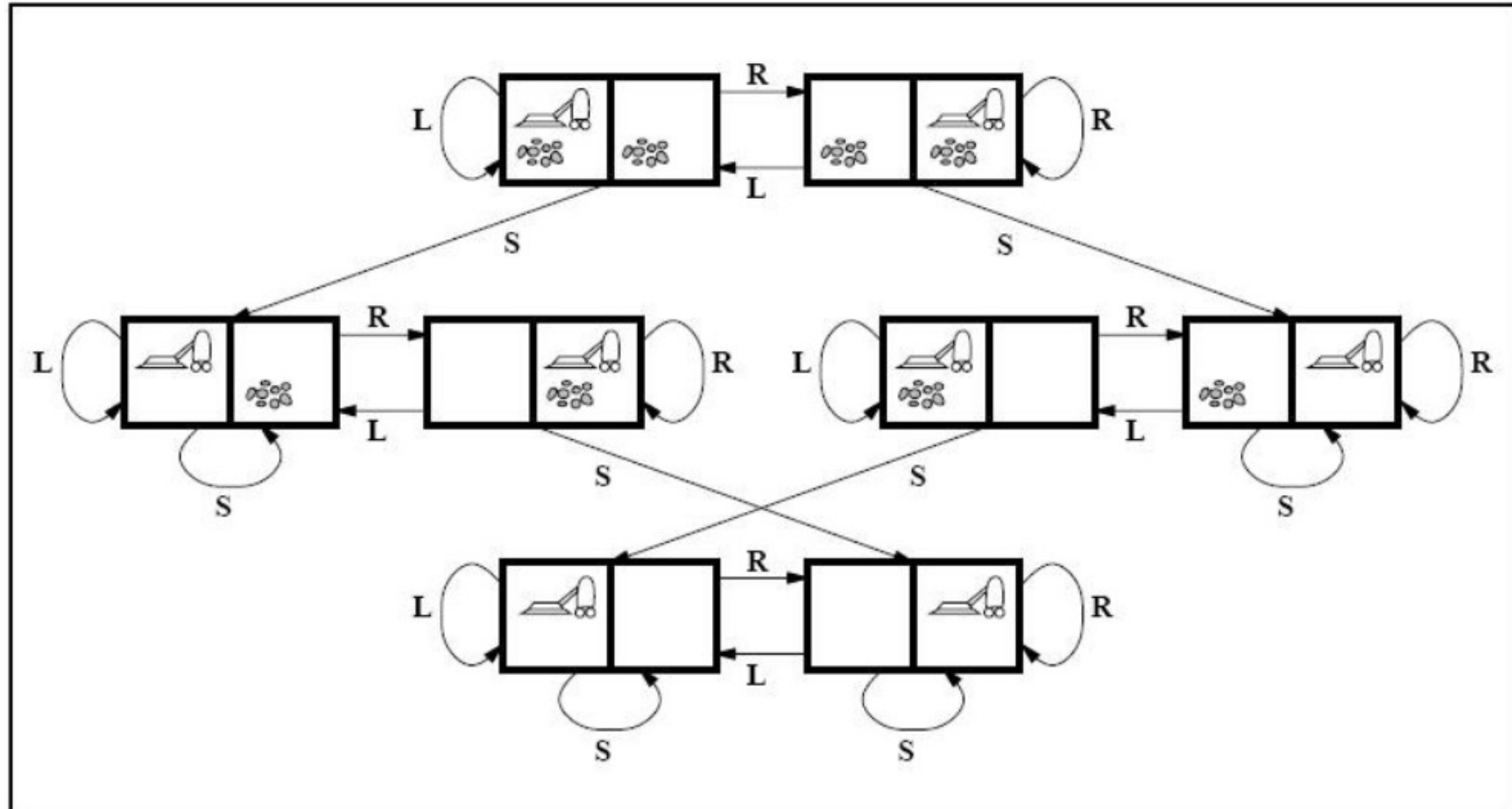
**Planning and Scheduling:  
Tutorial 1**

# Question

- Which of states below are valid states?
  - *Painted(LeftWall)*
  - *At(x)*
  - $\neg$ *Clean(Cotton)*
  - *Rainy(Tomorrow) or Cloudy(Tomorrow)*
  - *Passed(DueDate(307A3))*
  - *Hold(Banana) and At(RoomC)*



# PDDL in Vacuum Cleaner's World

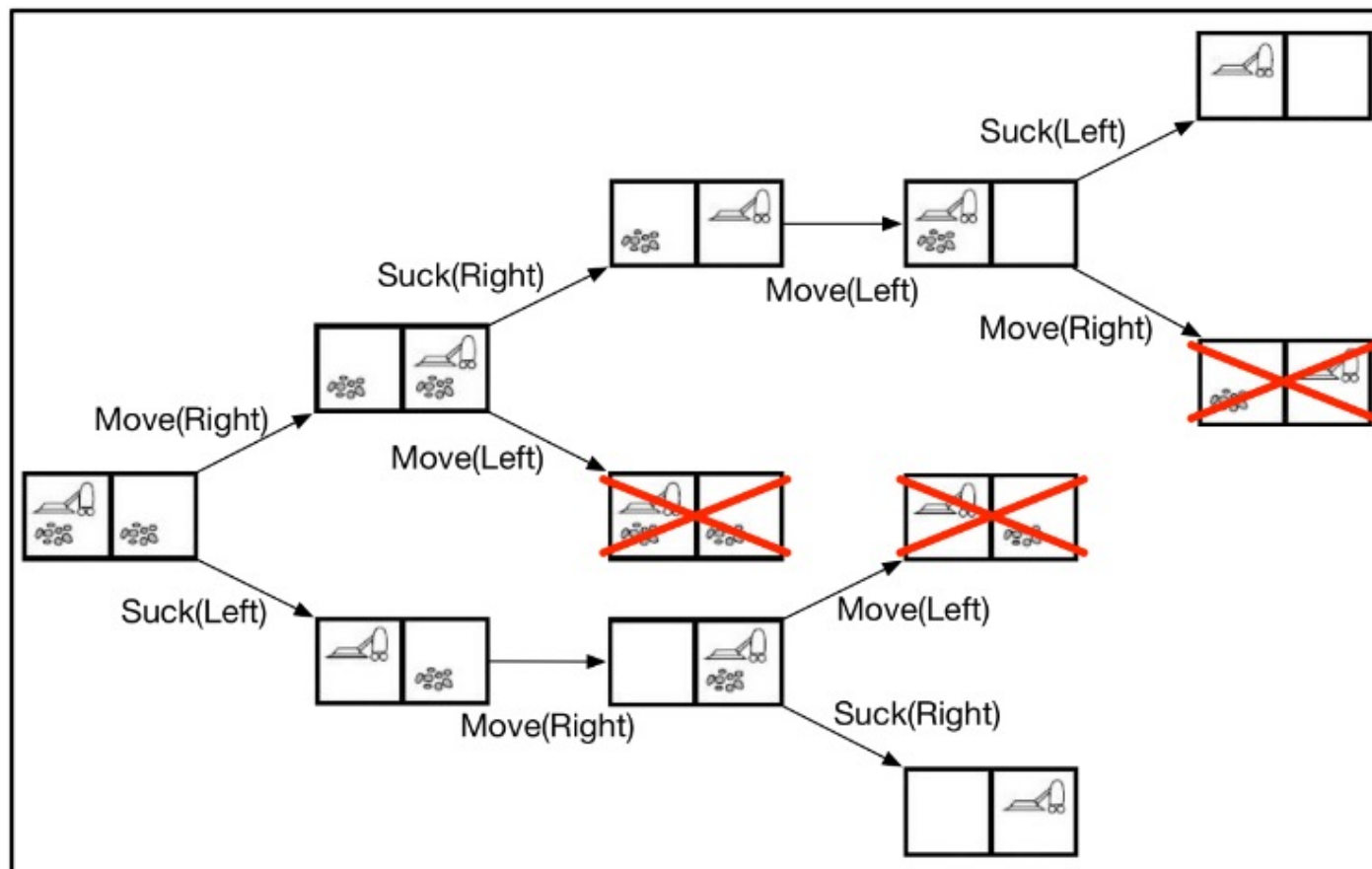


# Update State with Action

- Delete list  $\text{DEL}(a)$ : remove the fluents that appear as negative literals in the action's effects
- Add list  $\text{ADD}(a)$ : add the fluents that are positive literals in the action's effects
- $s' = \text{RESULT}(s, a) = (s - \text{DEL}(a)) \cup \text{ADD}(a)$
- Example in the vacuum cleaner's world
  - $s_1 = \text{At}(\text{Left}), a_1 = \text{MoveRight}()$ 
    - $\text{EFFECT}(a_1) = \text{At}(\text{Right}) \wedge \neg \text{At}(\text{Left})$
    - $s_1 - \text{DEL}(a_1) = \emptyset$
    - $\text{RESULT}(s_1, a_1) = \emptyset \cup \text{ADD}(a_1) = \text{At}(\text{Right})$
  - $s_2 = \text{At}(\text{Right}), a_2 = \text{Suck}(\text{Right})$ 
    - $\text{EFFECT}(a_2) = \text{Clean}(\text{Right})$
    - $s_2 - \text{DEL}(a_2) = \text{At}(\text{Right})$
    - $\text{RESULT}(s_2, a_2) = \text{At}(\text{Right}) \cup \text{ADD}(a_2) = \text{At}(\text{Right}) \wedge \text{Clean}(\text{Right})$

# Planning Algorithms as State-Space Search

- Forward (progression) state-space search
  - Start with the **initial state**
  - Examine all the **applicable actions** for the current state
  - **Avoid loop** – never go back to previous states
  - Until reach a **goal state**





# Question



- The “have cake and eat cake” planning problem:

*Init(Have(Cake, A))*

*Goal(Eaten(Cake, A)  $\wedge$  Eaten(Cake, B))*

*Action(Eat(Cake, x),*

*PRECOND : Have(Cake, x)*

*EFFECT :  $\neg$ Have(Cake, x)  $\wedge$  Eaten(Cake, x))*

*Action(Bake(Cake, x),*

*PRECOND :  $\neg$ Have(Cake, x)*

*EFFECT : Have(Cake, x))*

- For the initial state, list all the applicable actions and the resultant state for each of them.
- List three different plans (sequences of actions from the initial state to the goal state).

# JSS: an Example (Table)

- A solution is a schedule that processes these jobs with the machines (Gantt chart)

Job	Operation	Machine	ProcTime
1	AddEngine1	EngineHoist	30
	AddWheels1	WheelStation	30
	Inspect1	Inspector	10
2	AddEngine2	EngineHoist	60
	AddWheels2	WheelStation	15
	Inspect2	Inspector	10



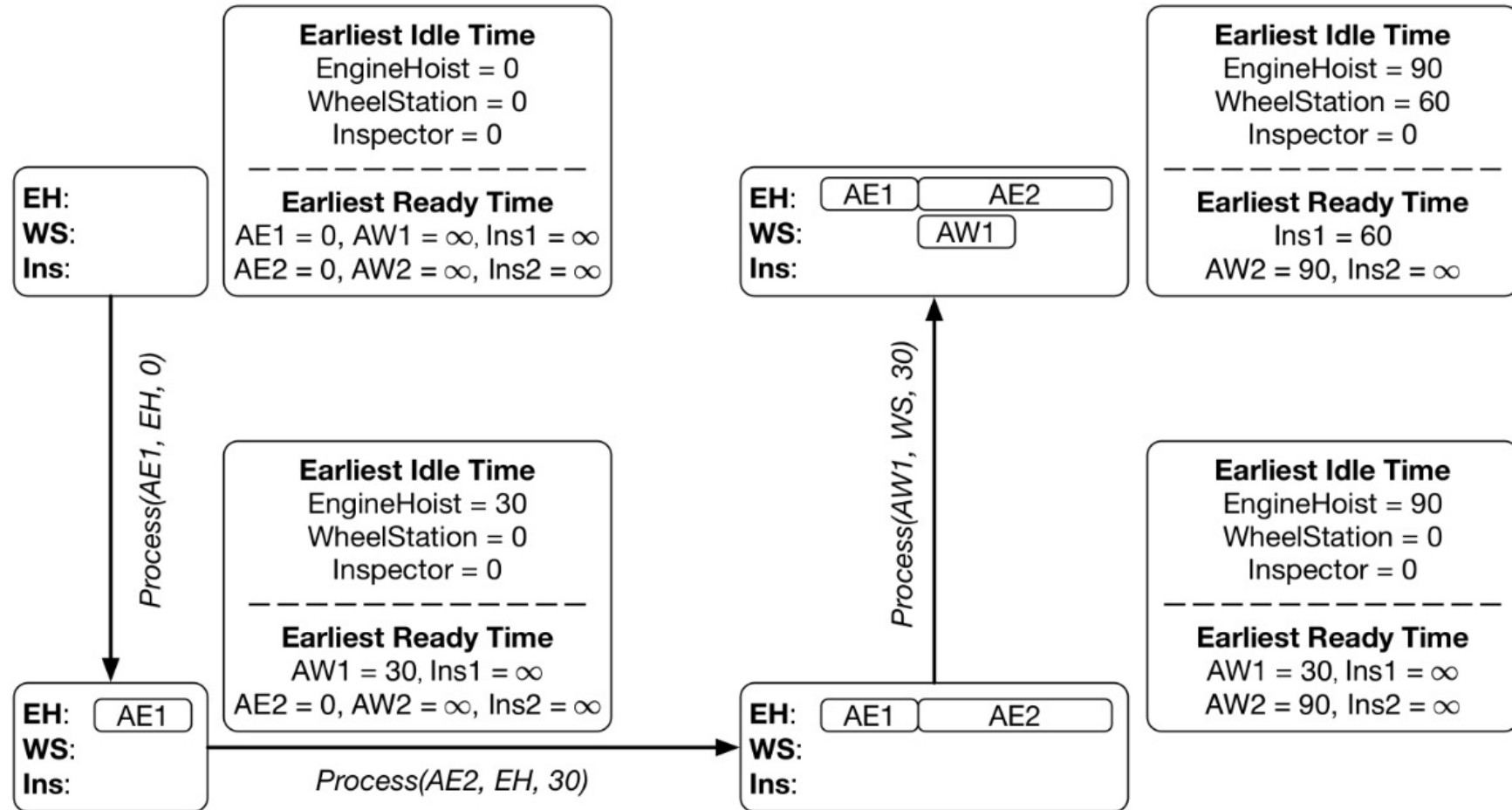
# Search for Schedules

- **Initial state**
  - Empty schedule,  $t=0$ , all operations unprocessed
  - The first operation of each job is ready at time 0, all the other operations are not ready
  - All machines are idle at time 0
- **Goal state**: all operations processed
- **Actions**:  $\text{Process}(o, m, t)$ 
  - Start processing operation  $o$  with machine  $m$  at time  $t$
  - Precondition:
    - $o$  unprocessed, and is ready at time  $t$
    - $m$  is idle at  $t$
  - Effect:
    - $o$  processed
    - $\text{next}(o)$  (if exists) is ready at time  $t + \text{ProcTime}(o)$ , and  $m$  is idle at  $t + \text{ProcTime}(o)$
- How to decide  $t$ ?

# Deciding Starting Time of Action

- **Non-delay**: start the action as soon as possible
  - Operation **earliest ready time**
  - Machine **earliest idle time**
  - **Earliest starting time**: the **later** between the above two
- Find operation earliest ready time
  - Initial: 0 for the first operation, and infinity for others
  - When  $\text{Process}(o,m,t)$  is scheduled, then the earliest ready time of  $\text{next}(o)$  becomes  $t + \text{ProcTime}(o)$
- Find machine earliest idle time
  - Initial: 0
  - When  $\text{Process}(o,m,t)$  is scheduled, then the earliest idle time of machine  $m$  becomes  $t + \text{ProcTime}(o)$

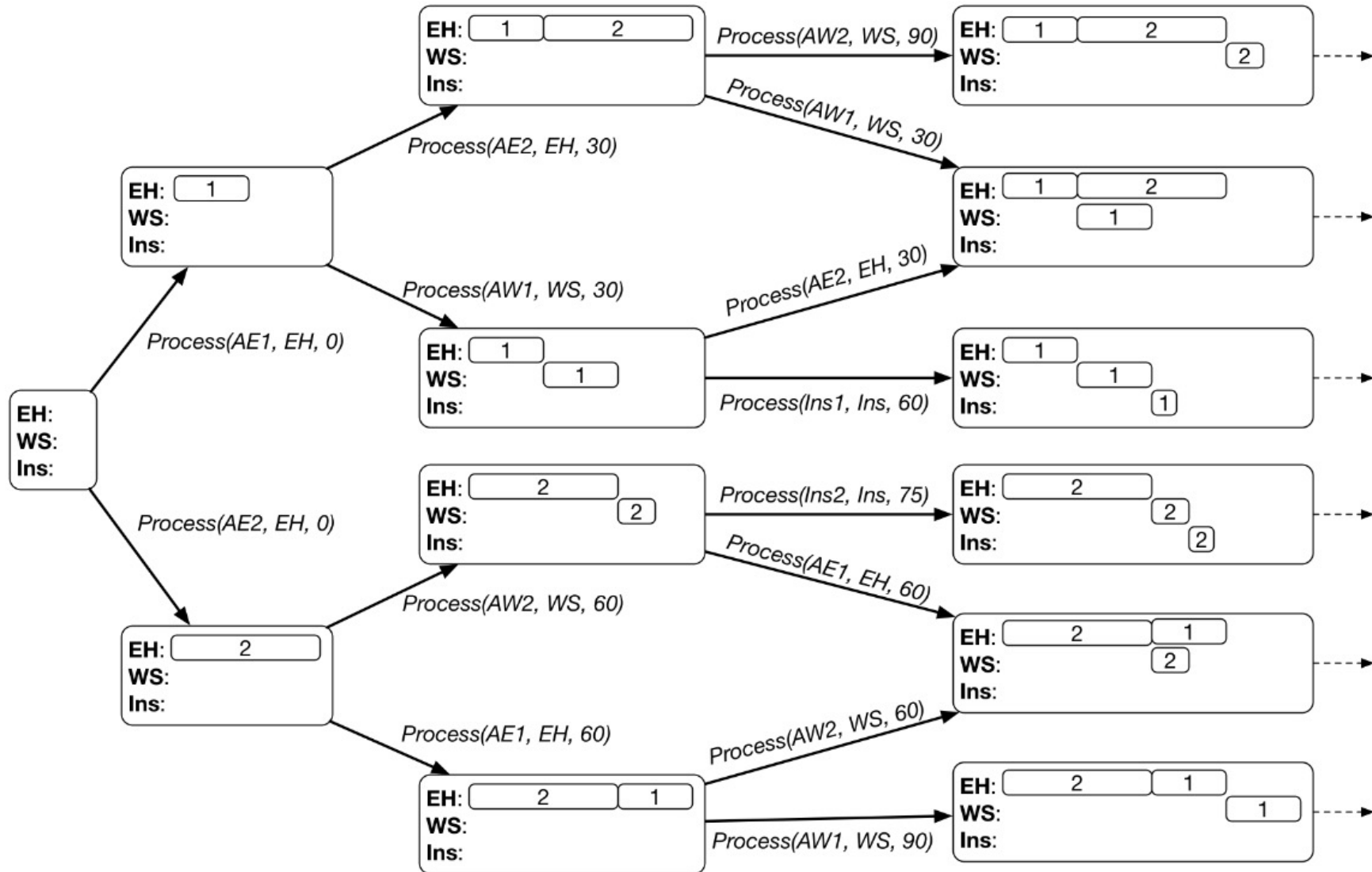
# Update Earliest Ready and Idle Time



# Forward State-Space Search

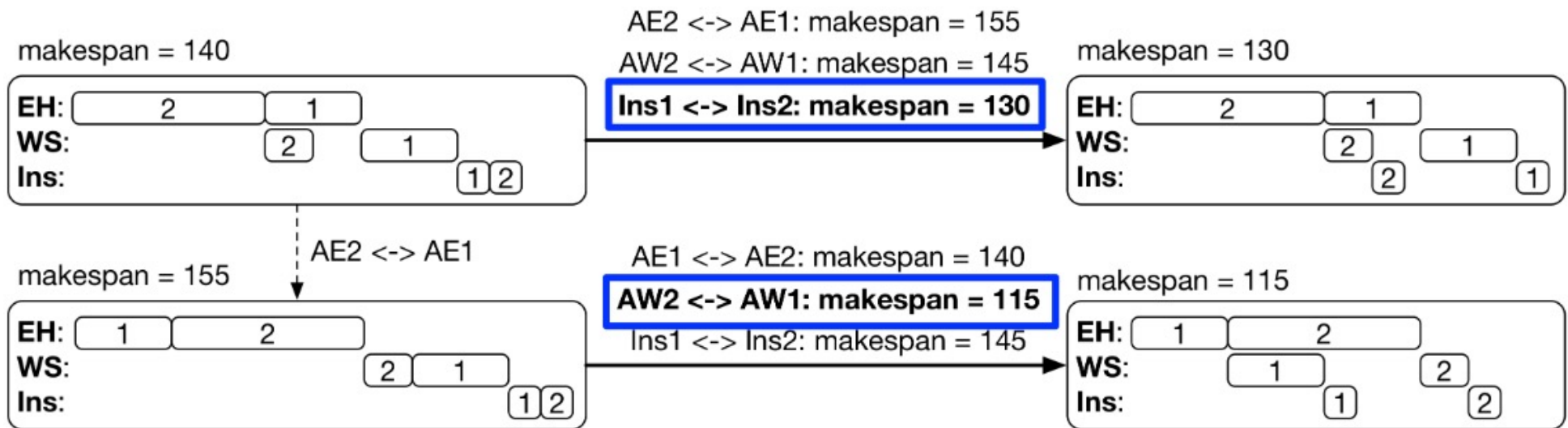
- Start from the initial state
  - Empty schedule,  $t=0$ , all operations unprocessed
  - The first operation of each job is **ready** at time 0, all the other operations are not ready
  - All machines are **idle** at time 0
- Examine all the **applicable** actions  $\text{Process}(o,m,t)$ 
  - Enumerate each unprocessed operation  $o$  and its machine  $m$
  - Calculate the earliest starting time  $t$
  - Applicable if  $t < \infty$
- All the leaf nodes are goal states (all operations processed)
- Each schedule is a path from the root node to a leaf node

# Forward State-Space Search



# Local Search (Hill Climbing)

- Step 1. Random generate a scheduling  $s$ ;
- Step 2. Examine all the neighbors in the neighborhood of  $s$ , and select the best neighbor  $s'$ ;
- Step 3. If  $s'$  is better than  $s$ , set  $s \leftarrow s'$ , and go to Step 2. Otherwise return  $s$ .



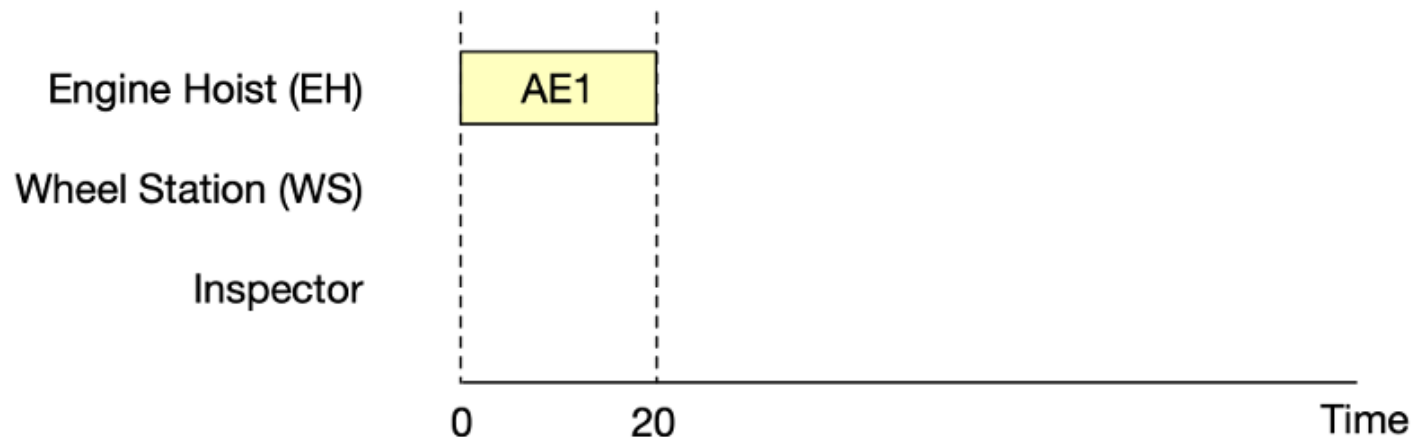
- Jump out of local optima: simulated annealing, genetic algorithms, ...

# Question

- For the car manufacturing scheduling problem, consider 3 jobs as summarized below:

	Arrival Time	Processing Time		
		AddEngine ( <i>AE</i> )	AddWheels ( <i>AW</i> )	Inspect ( <i>Ins</i> )
Job 1	0	20	30	15
Job 2	0	45	20	20
Job 3	0	25	30	10

- Given the partial schedule below:



- List all applicable actions in this state, formatted as (Operation, Machine, StartTime)