

### **NWEN 243**

**Networked Applications** 

Transport layer and application layer





# Topic

- TCP flow control
- TCP congestion control
- The Application Layer



#### Fast Retransmit

- Time-out period often relatively long:
  - long delay before resending lost packet
- Detect lost segments via duplicate ACKs.
  - Sender often sends many segments back-to-back
  - If segment is lost, there will likely be many duplicate ACKs.

- If sender receives 3 ACKs for the same data, it supposes that segment after ACKed data was lost:
  - <u>fast retransmit</u>: resend segment before timer expires



#### **TCP Flow Control**

#### - flow control

sender won't overrun Receiver's buffers by transmitting too much, too fast

**RcvBuffer** = size or TCP Receive Buffer



receiver: explicitly informs sender of (dynamically changing) amount of free buffer space

> RcvWindow field in TCP segment

sender: keeps the amount of transmitted, unACKed data less than most recently received RcvWindow



# Receive Window (Example)





## Quick exercise



- A. ack=2000, win=1000
- B. ack=2000, win=2000
- C. ack=4000, win=1000
- D. ack=4000, win=2000

### **Principles of Congestion Control**

#### Congestion:

- Informally: "too many sources sending too much data too fast for *network* to handle"
- Different from flow control!
- Manifestations:
  - lost packets (buffer overflow at routers)
  - long delays (queuing in router buffers)





### Causes/costs of congestion

- Four senders/receivers
- Multi-hop paths
- Timeout/retransmit





#### Extra "cost" of congestion:

when packet dropped, any upstream transmission capacity used for that packet was wasted!



# **TCP Slow Start**

- Probe a link's capacity...
- For every round of successful transmission, the congestion window size is doubled.
- First timeout establishes our first guess at the link's capacity.
- This topic is covered in detail in NWEN302.





## Quick exercise

- Which of the following are correct?
  - A. In the slow start algorithm, the congestion window increases additively until congestion is detected.
  - B. If congestion is detected by timeout, a new congestion avoidance phase will be started.
  - C. Congestion control aims to regulate the amount of data a source can send before receiving any acknowledgement.
  - D. None of the above

# Quick exercise on TCP

- A.TCP supports reliable transport: between sending and receiving process
- B.TCP supports flow control: sender won't overwhelm receiver
- C.TCP supports congestion control: throttle sender when network overloaded
- D.TCP provides timing, minimum bandwidth guarantees

# Quick exercise on UDP

- A. Best effort data transfer between sending and receiving process
- B. Does not require connection setup
- C. Does not check data integrity
- D. Does not provide flow control, congestion control, timing, or bandwidth guarantees
- E. Provide maximum freedom to application designers



# **Application Layer**

- Application characteristics
- Communication architecture
- HTTP



## Application-layer protocols

- Just as the Datalink, Network and Transport layers had protocols to communicate
- So too does the Application layer. However these are defined on a per application basis.
- Examples, HTTP, SMTP, Jabber/XMPP etc. There are many protocols, and no one can know them all (many are proprietary)

# **Communication Architectures**



# Application Protocol Characteristics

#### Data loss

- some apps (e.g., audio) can tolerate some loss
- other apps (e.g., file transfer, telnet) require 100% reliable data transfer

#### Timing

 some apps (e.g., Internet telephony, interactive games) require low delay to be "effective"

#### Bandwidth

some apps (e.g., multimedia)
 require minimum amount
 of bandwidth to be
 "effective"
other apps ("elastic apps")
 make use of whatever
 bandwidth they get



### Quick exercise

- Suppose that you are building an interactive computer gaming application, what communication requirements should the gaming protocol support?
  - A. minimum bandwidth
  - B. no data loss
  - C. maximum delay
  - D. connection-oriented delivery

# The Web: the HTTP protocol

- Client-server model:
  - client: browser that requests, receives, "renders" Web objects
  - server: Web server sends objects in response to requests





# A Little Historical Context

- The `WEB' is composed of a set of linked HTML documents.
- The WWW is a networl
- The HyperText Transfer is an application level pro 'web' documents over the



## Growth of a Protocol (HTTP)

- Started with HTTP/0.9 (1990)
  - raw data transfer over the Internet (simple)
  - client issues GET Request
  - server replies with bytes comprising the HTML

Client request line (GET)

**GET /somedir/page.html** (extra carriage return, line feed)

blah blah blah blah...

Server data, e.g., requested html file

# Growth of a Protocol (HTTP/I.0)

- Next HTTP/1.0 (1996)
  - sophisticated requests {GET, HEAD and POST}
  - server response {headers + data}
  - better client-server coordination/interaction.
  - coordinates with other protocols {SMTP, NNTP, ...}
  - MIME (Multipurpose Internet Mail Extensions)