

## NWEN 243

Networked Applications

Transport layer and application layer



## Topic

- TCP flow control
- TCP congestion control
- The Application Layer

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## Fast Retransmit

- Time-out period often relatively long:
  - long delay before resending lost packet
- Detect lost segments via duplicate ACKs.
  - Sender often sends many segments back-to-back
  - If segment is lost, there will likely be many duplicate ACKs.
- If sender receives **3 ACKs** for the same data, it supposes that segment after ACKed data was lost:
  - **fast retransmit**: resend segment before timer expires

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## TCP Flow Control

**flow control**  
sender won't overrun Receiver's buffers by transmitting too much, too fast

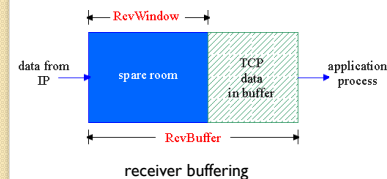
**receiver**: explicitly informs sender of (dynamically changing) amount of **free buffer space**

- **RcvWindow field** in TCP segment

**RcvBuffer** = size of TCP Receive Buffer

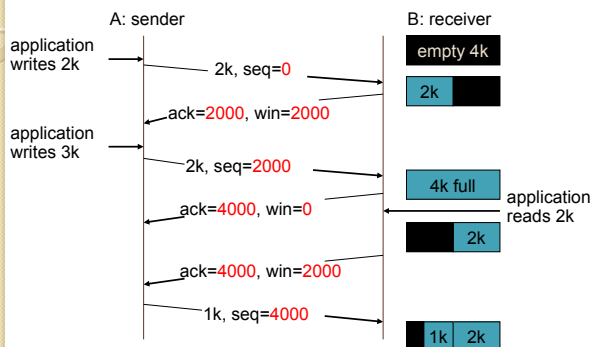
**RcvWindow** = amount of spare room in Buffer

**sender**: keeps the amount of **transmitted, unACKed** data less than most recently received **RcvWindow**



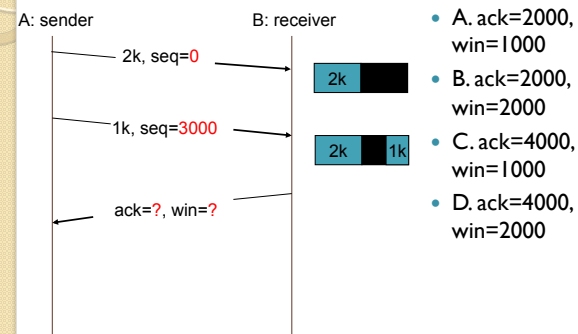
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## Receive Window (Example)



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## Quick exercise



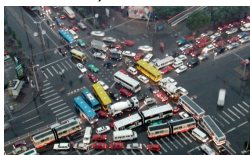
- A. ack=2000, win=1000
- B. ack=2000, win=2000
- C. ack=4000, win=1000
- D. ack=4000, win=2000

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## Principles of Congestion Control

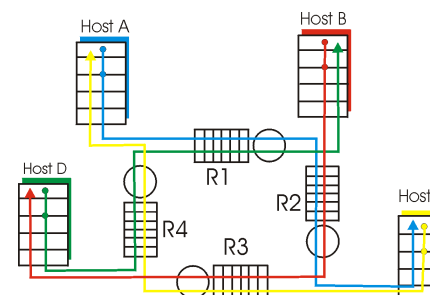
### Congestion:

- Informally: “too many sources sending too much data too fast for network to handle”
- Different from flow control!
- Manifestations:
  - lost packets (buffer overflow at routers)
  - long delays (queuing in router buffers)

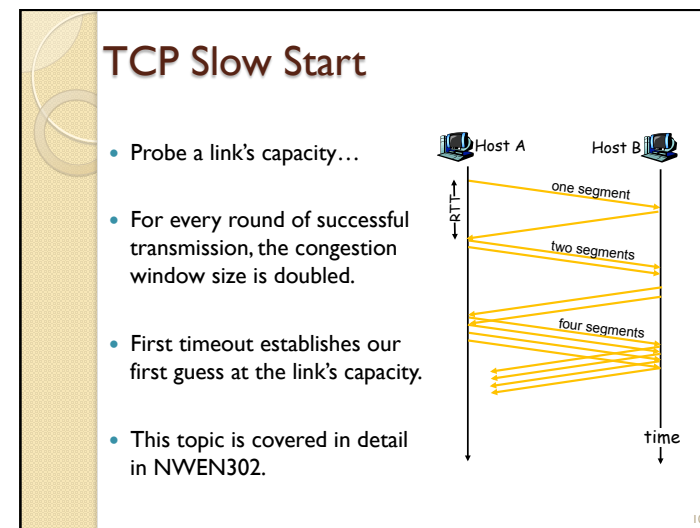
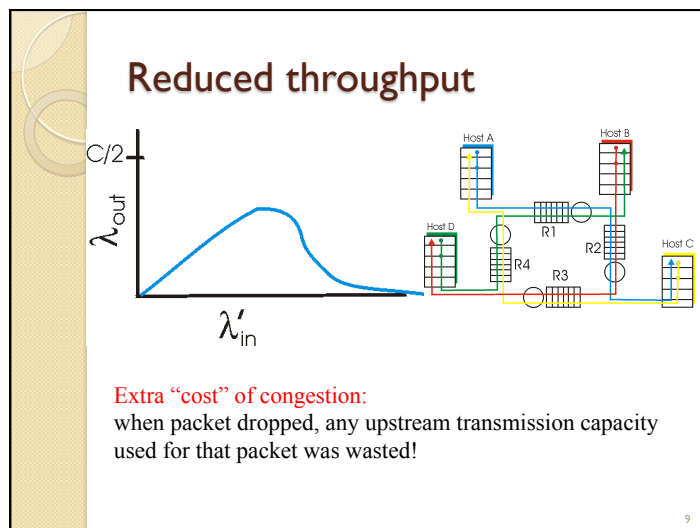


## Causes/costs of congestion

- Four senders/receivers
- Multi-hop paths
- Timeout/retransmit



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- ### Quick exercise
- Which of the following are correct?
    - A. In the slow start algorithm, the congestion window increases additively until congestion is detected.
    - B. If congestion is detected by timeout, a new congestion avoidance phase will be started.
    - C. Congestion control aims to regulate the amount of data a source can send before receiving any acknowledgement.
    - D. None of the above

- ### Quick exercise on TCP
- A. TCP supports reliable transport: between sending and receiving process
  - B. TCP supports flow control: sender won't overwhelm receiver
  - C. TCP supports congestion control: throttle sender when network overloaded
  - D. TCP provides timing, minimum bandwidth guarantees

## Quick exercise on UDP

- A. Best effort data transfer between sending and receiving process
- B. Does not require connection setup
- C. Does not check data integrity
- D. Does not provide flow control, congestion control, timing, or bandwidth guarantees
- E. Provide maximum freedom to application designers

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## Application Layer

- Application characteristics
- Communication architecture
- HTTP

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## Application-layer protocols

- Just as the Datalink, Network and Transport layers had protocols to communicate
- So too does the Application layer. However these are **defined on a per application basis**.
- Examples, HTTP, SMTP, Jabber/XMPP etc. There are many protocols, and no one can know them all (many are proprietary)

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## Communication Architectures

- Messaging



- Request-response



- Do-operation



- Streaming



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## Application Protocol Characteristics

### Data loss

- some apps (e.g., audio) can tolerate some loss
- other apps (e.g., file transfer, telnet) require 100% reliable data transfer

### Timing

- some apps (e.g., Internet telephony, interactive games) require low delay to be “effective”

### Bandwidth

- some apps (e.g., multimedia) require minimum amount of bandwidth to be “effective”
- other apps (“elastic apps”) make use of whatever bandwidth they get

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## Quick exercise

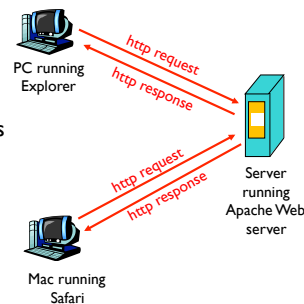
- Suppose that you are building an **interactive computer gaming application**, what communication requirements should the gaming protocol support?
  - A. minimum bandwidth
  - B. no data loss
  - C. maximum delay
  - D. connection-oriented delivery

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## The Web: the HTTP protocol

### Client-server model:

- **client**: browser that requests, receives, “renders” Web objects
- **server**: Web server sends objects in response to requests



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## A Little Historical Context

- The ‘**WEB**’ is composed of a set of **linked HTML documents**.
- The **WWW** is a network
- The **HyperText Transfer** is an **application** level protocol for transferring ‘web’ documents over the network



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## Growth of a Protocol (HTTP)

- Started with HTTP/0.9 (1990)
  - raw data transfer over the Internet (simple)
  - client issues **GET Request**
  - server replies with bytes comprising the HTML

Client  
request line  
(GET) → **GET /somedir/page.html**  
(extra carriage return, line feed)

Server data,  
e.g.,  
requested  
html file ←  
blah blah blah blah...

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## Growth of a Protocol (HTTP/1.0)

- Next HTTP/1.0 (1996)
  - sophisticated requests {GET, HEAD and POST}
  - server response {headers + data}
  - better client-server coordination/interaction.
  - coordinates with other protocols {SMTP, NNTP, ...}
  - MIME (Multipurpose Internet Mail Extensions)

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