

ENGR 101

Engineering Technology

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Victoria
UNIVERSITY OF WELLINGTON
*Te Whare Wānanga
o te Ūpoko o te Ika a Māui*



CAPITAL CITY UNIVERSITY

Week 3 Lecture 4b

- Main topics
 - Introduction to Engineering Technology
 - Number system
 - Logic Gates
 - Boolean Algebra
- Course web page:
https://ecs.wgtn.ac.nz/Courses/XMUT101_2021T1/
- kerese@ecs.vuw.ac.nz

Binary Addition

- Binary addition is similar to decimal addition.
- Adding 2 binary numbers:

$$1\ 1\ 1\ 1\ 0\ 1 + 1\ 0\ 1\ 1\ 1$$

Binary Addition

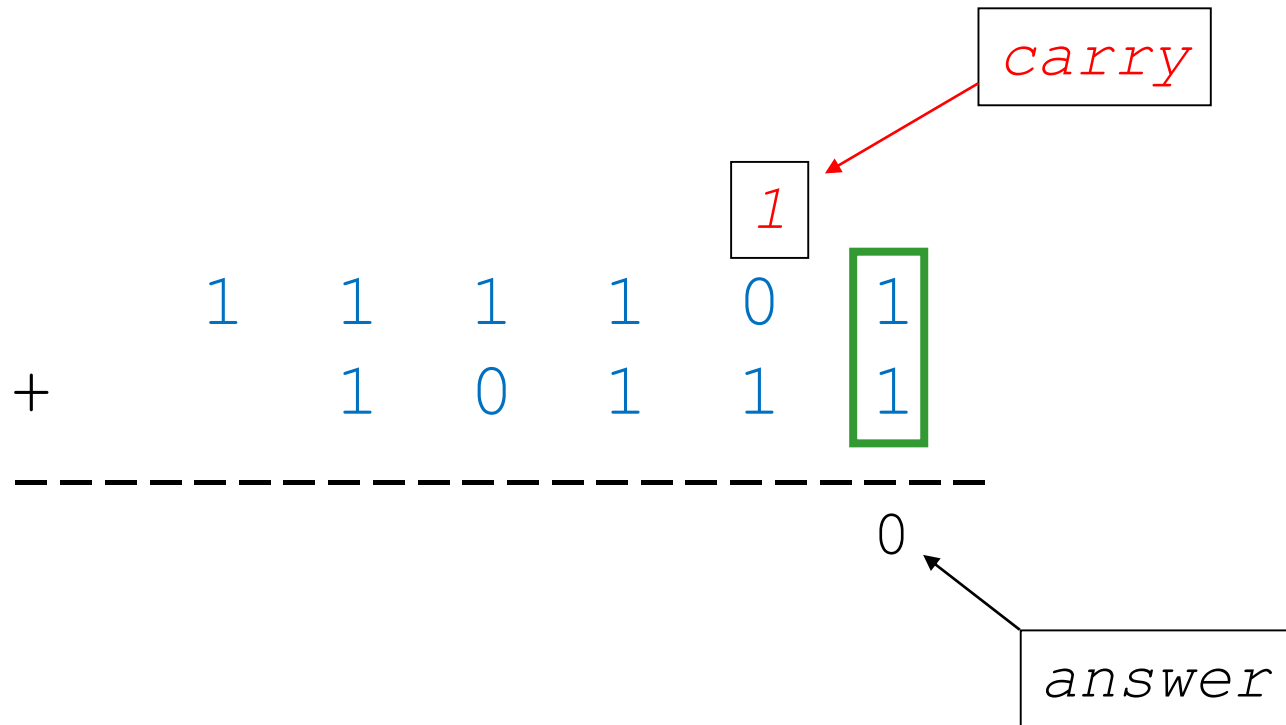
- Adding 2 binary numbers: $111101 + 10111$

$$\begin{array}{r} \\ + \\ \hline \end{array}$$

The diagram shows the binary addition of 111101 and 10111. The numbers are aligned by their least significant bits. A dashed line is drawn below the numbers, and a '0' is written below it, indicating a carry into the next column. The rightmost column (the 6th bit from the left) contains two '1's, which are highlighted with a green box, representing the sum of 1 + 1 = 10 in binary, where the '0' is the carry.

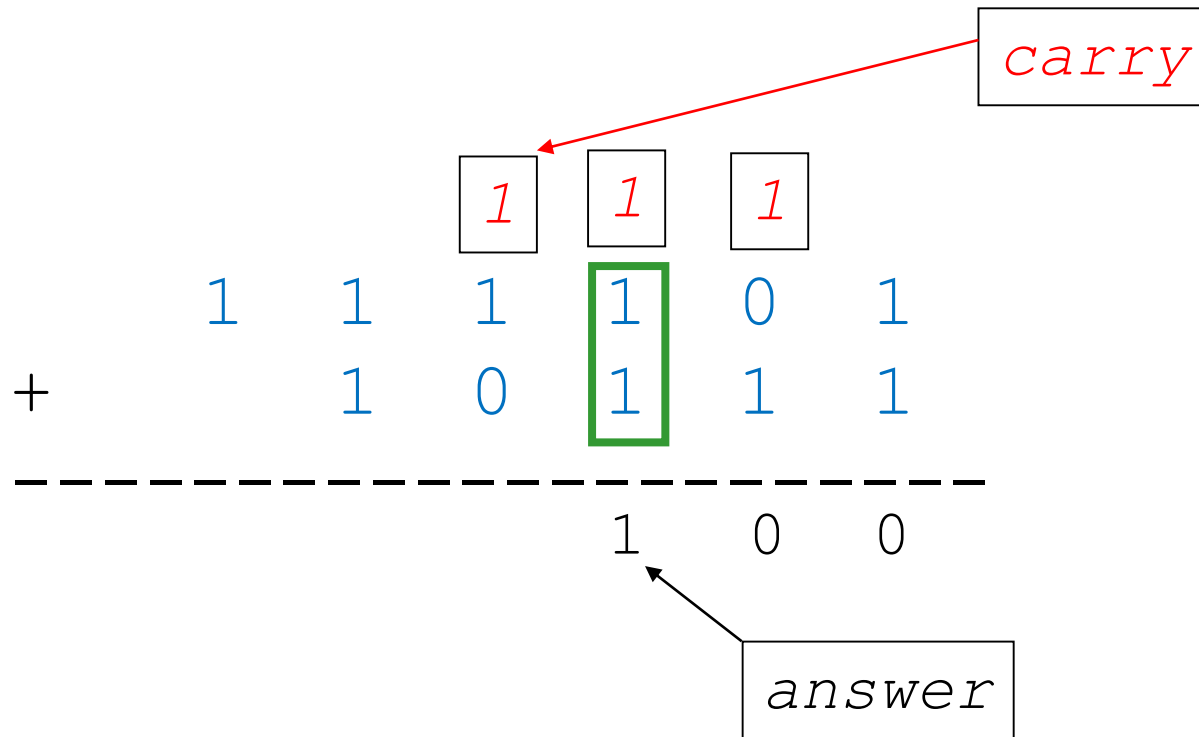
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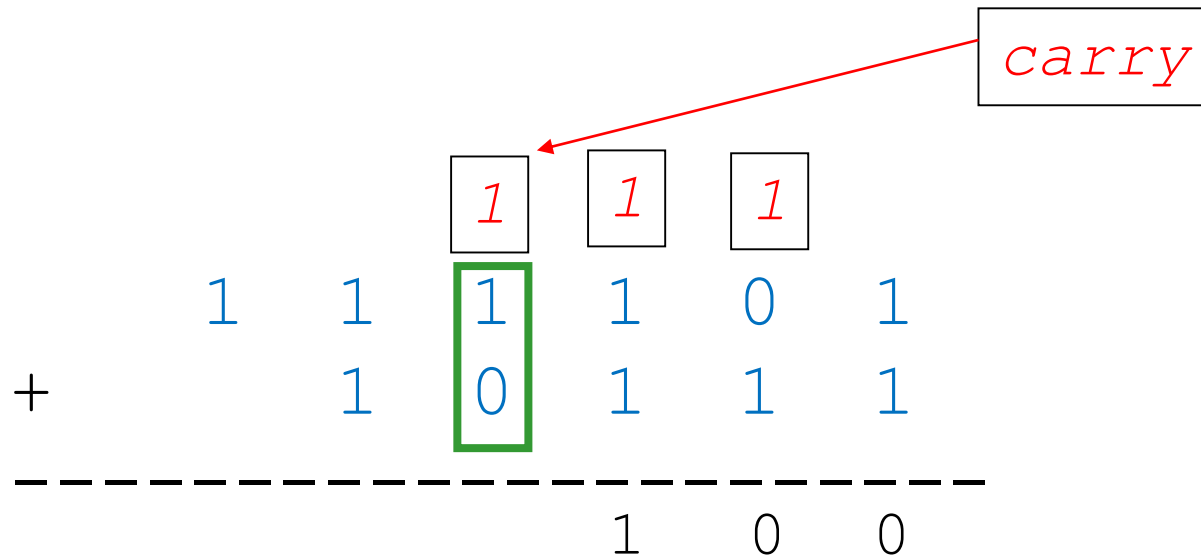
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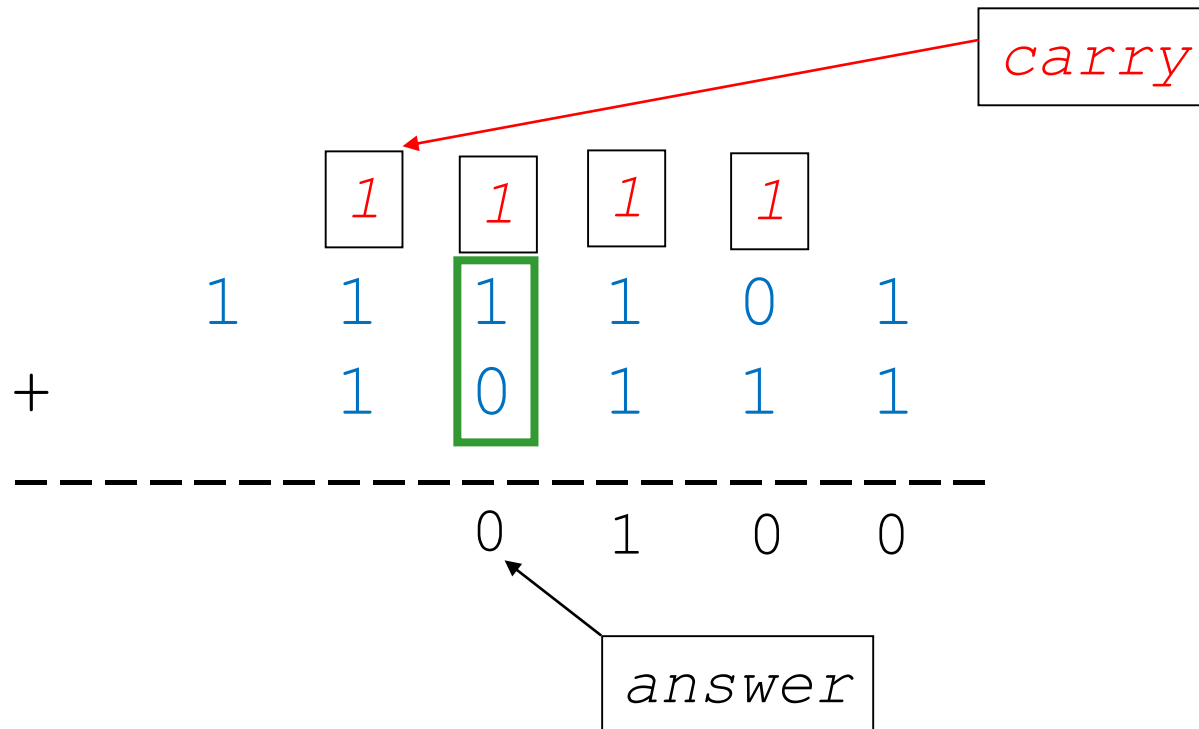
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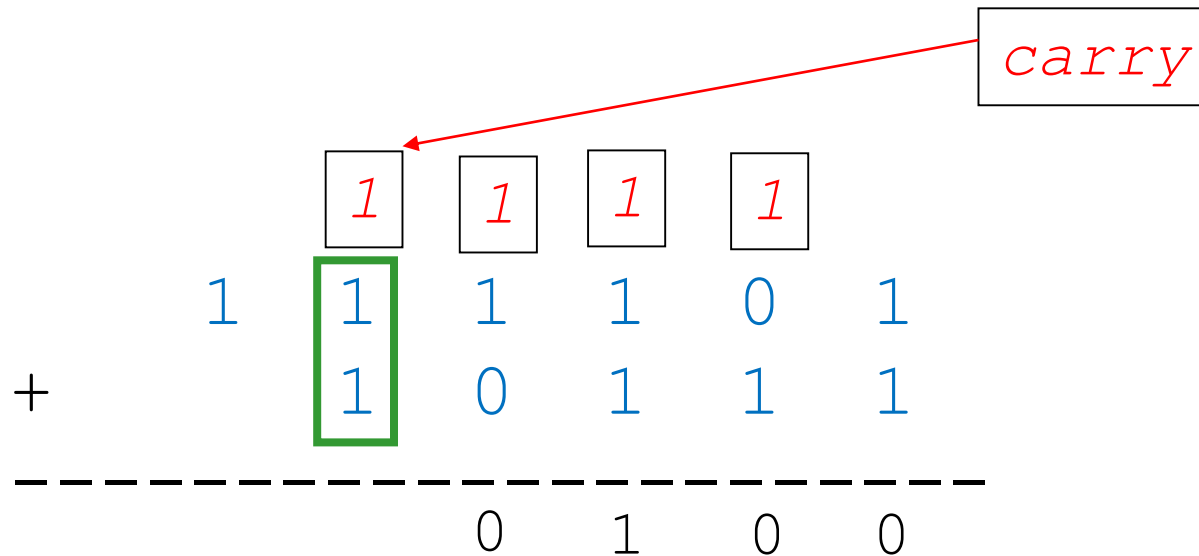
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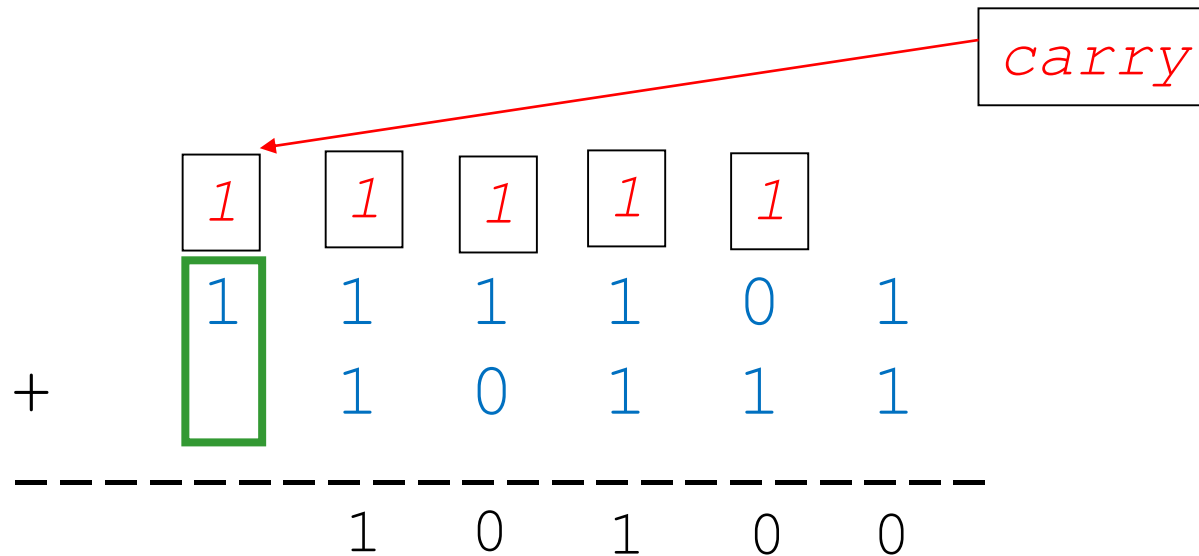
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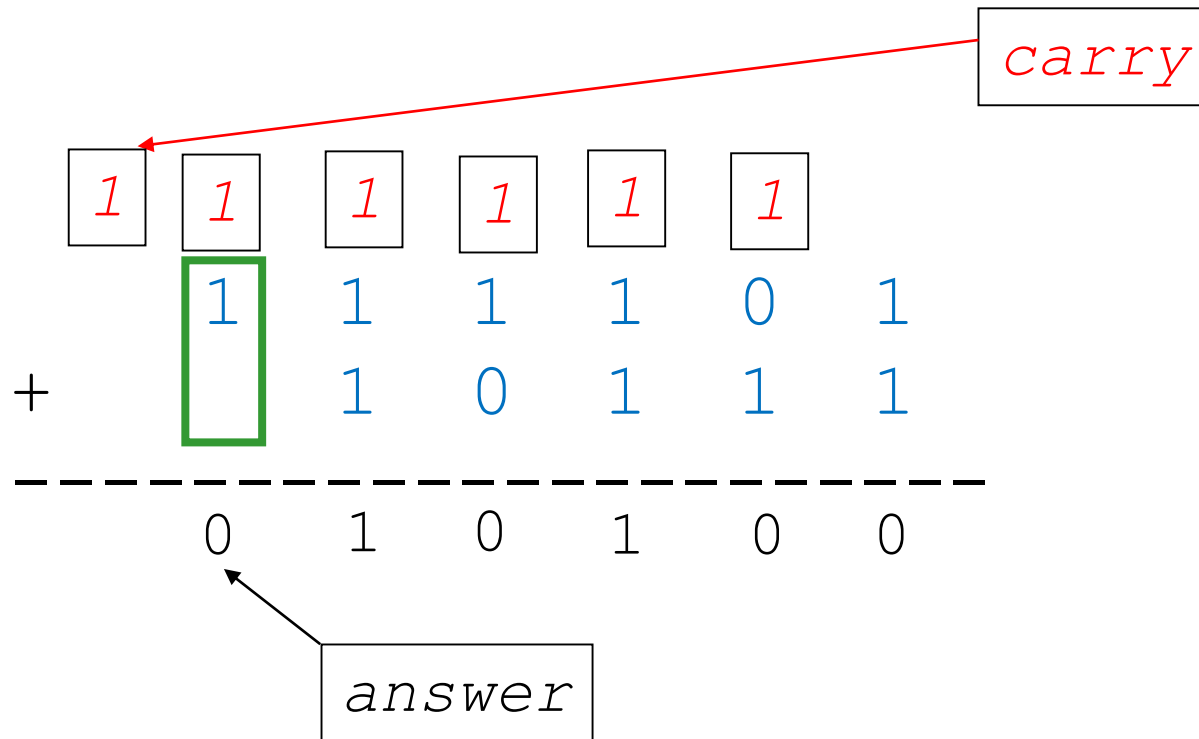
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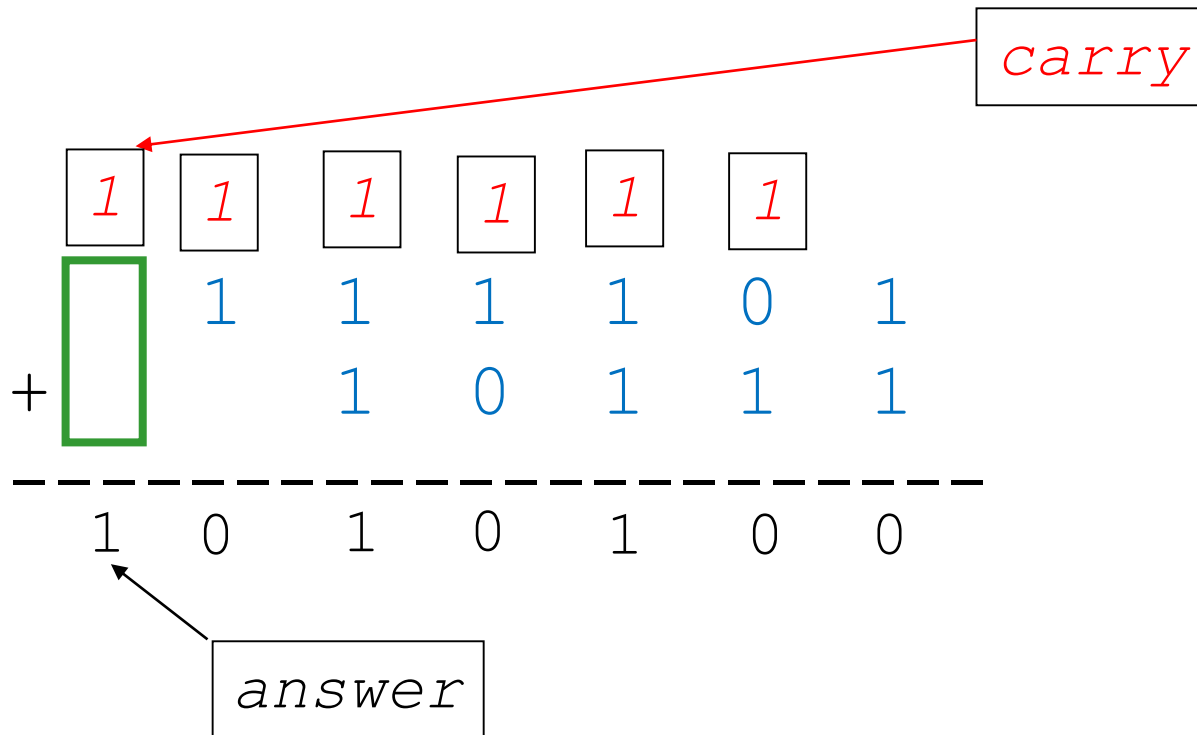
Binary Addition

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$$\begin{array}{rcccccc} & \boxed{1} & \boxed{1} & \boxed{1} & \boxed{1} & \boxed{1} & & \\ & & 1 & 1 & 1 & 1 & 0 & 1 \\ + & & & 1 & 0 & 1 & 1 & 1 \\ \hline & 1 & 0 & 1 & 0 & 1 & 0 & 0 \end{array}$$



Is this correct?

Binary Addition

- Adding 2 binary numbers: $111101 + 10111$

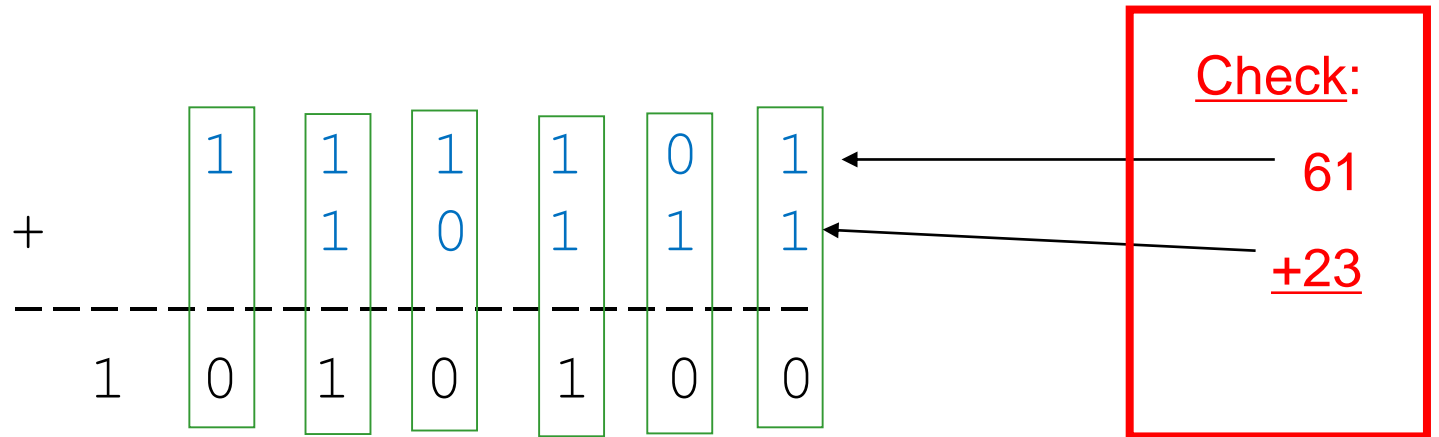
$$\begin{aligned} & (1 \times 2^5) + (1 \times 2^4) + (1 \times 2^3) + (1 \times 2^2) + (0 \times 2^1) + (1 \times 2^0) \\ &= (1 \times 32) + (1 \times 16) + (1 \times 8) + (1 \times 4) + (0 \times 2) + (1 \times 1) \\ &= 32 + 16 + 8 + 4 + 0 + 1 \\ &= 48 + 13 \\ &= 61 \end{aligned}$$

Check:

+	1	1	1	1	0	1	←	61
	0	1	0	1	0	0		
1								

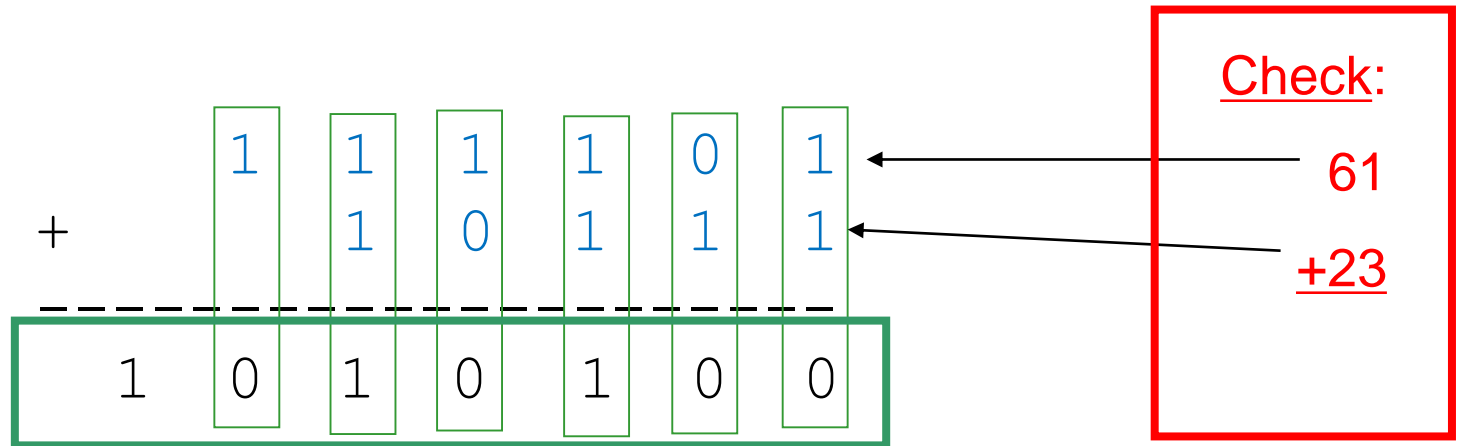
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Binary Addition

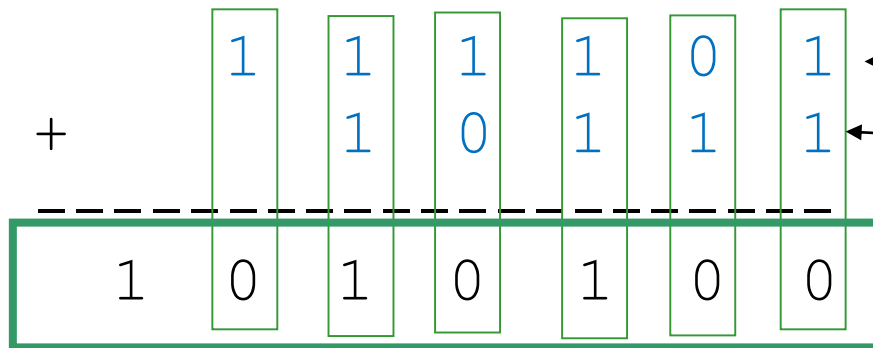
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Binary Addition

- Adding 2 binary numbers: $111101 + 10111$

$$\begin{aligned} & (1 \times 2^6) + (0 \times 2^5) + (1 \times 2^4) + (0 \times 2^3) + (1 \times 2^2) + (0 \times 2^1) + (0 \times 2^0) \\ &= (1 \times 64) + (0 \times 32) + (1 \times 16) + (0 \times 8) + (1 \times 4) + (0 \times 2) + (0 \times 1) \\ &= 64 + 0 + 16 + 0 + 4 + 0 + 0 \\ &= 64 + 20 \\ &= 84 \end{aligned}$$



Check:

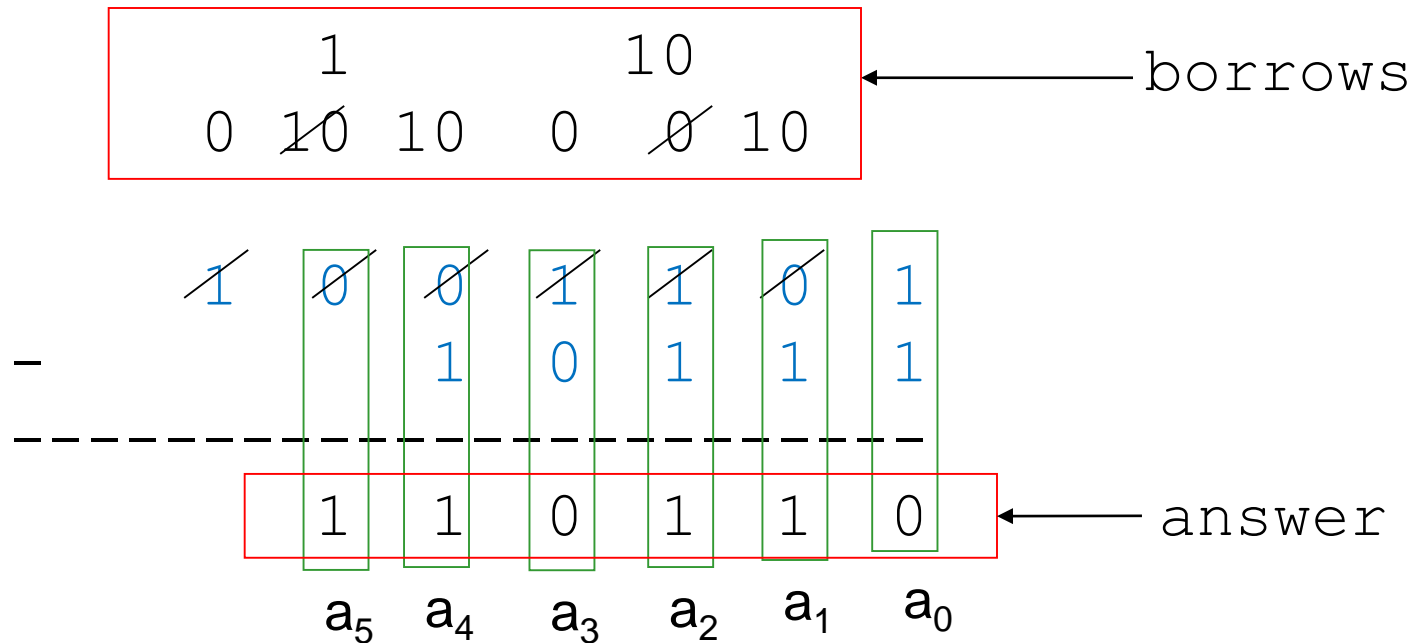
61

+23

84

Binary Subtraction

- We can also perform subtraction (with borrows in place of carries).
- Let's subtract $(10111)_2$ from $(1001101)_2$



How To Represent Signed Numbers?

- Plus (+) and minus (-) signs are used for decimal numbers: 25 (or +25), -16, etc.

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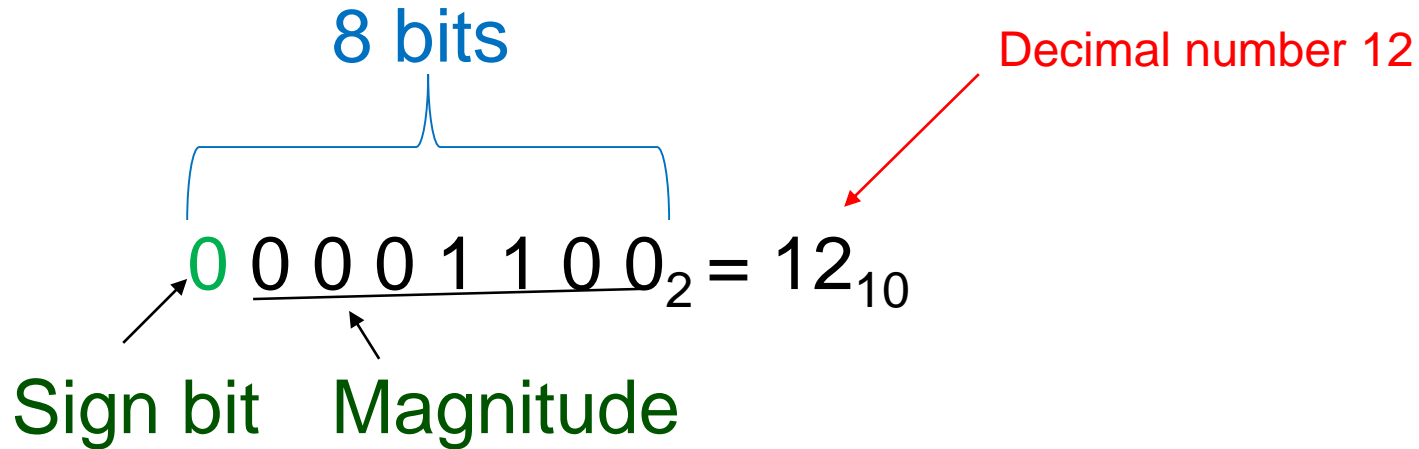
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25 (or +25), -16, etc.
- For computers, desirable to represent everything as *bits*. (Bit – Binary digit)
- Three types of signed binary number representations:
signed magnitude, 1's complement, 2's complement.

How To Represent Signed Numbers

- Plus and minus sign used for decimal numbers: 25 (or +25), -16, etc.
 - For computers, desirable to represent everything as *bits*.
 - Three types of signed binary number representations: signed magnitude, 1's complement, 2's complement.
- In each case: left-most bit indicates the sign; positive (0) or negative (1).

1) Signed magnitude numbers

Signed magnitude example:



1) Signed magnitude numbers

Consider *signed magnitude*:

$$\begin{array}{c} \nearrow \text{00001100}_2 = 12_{10} \\ \text{Sign bit} \quad \text{Magnitude} \end{array}$$

$$\begin{array}{c} \nearrow \text{10001100}_2 = -12_{10} \\ \text{Sign bit} \quad \text{Magnitude} \end{array}$$

2) One's Complement Representation

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- To find negative of 1's complement number take the 1's complement.

Sign bit Magnitude

$$\underline{00001100}_2 = 12_{10}$$

Sign bit Magnitude

$$\underline{11110011}_2 = -12_{10}$$

3) Two's Complement Representation

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- The two's complement of a binary number involves inverting all bits and adding 1.
 - 2's comp of 00110011 is 11001101
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- For a n-bit number N the 2's complement is:

$$(2^n - 1) - N + 1$$

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- To find negative of 2's complement number take the 2's complement.

$00001100_2 = 12_{10}$

Sign bit Magnitude

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Sign bit Magnitude

Two's Complement Shortcuts

Algorithm 1 – Simply complement each bit and then add 1 to the result.

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$$\begin{array}{r} N = 01100101 \\ \begin{array}{cccccccc} \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow \\ 1 & 0 & 0 & 1 & 1 & 0 & 1 & 0 \end{array} \\ + \qquad \qquad \qquad 1 \\ \hline \boxed{10011011} \end{array}$$

Invert all bits in N

Add 1

Two's Complement Shortcuts

- Algorithm 1 – Simply complement each bit and then add 1 to the result.
 - Find the 2's complement of $(01100101)_2$ and of **its 2's complement...**

$$\begin{array}{r} N = 01100101 \\ \quad 10011010 \\ + \quad \quad 1 \\ \hline 10011011 \end{array}$$

$$\begin{array}{r} [N] = 10011011 \\ \quad 01100100 \\ + \quad \quad 1 \\ \hline 01100101 \end{array}$$

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- Algorithm 1 – Simply complement each bit and then add 1 to the result.
 - Finding the 2's complement of $(01100101)_2$ and of its 2's complement...

N =	01100101	[N] =	10011011
	10011010		01100100
+	1	+	1
-----		-----	
	10011011		01100101

- Algorithm 2 – Start with the least significant bit, copy all of the bits up to and including the first 1 bit and then complementing the remaining bits.

Two's Complement Shortcuts

- Algorithm 1 – Simply complement each bit and then add 1 to the result.
 - Finding the 2's complement of $(01100101)_2$ and of its 2's complement...

$$\begin{array}{r} N = 01100101 \\ 10011010 \\ + \quad \quad 1 \\ \hline 10011011 \end{array} \qquad \begin{array}{r} [N] = 10011011 \\ 01100100 \\ + \quad \quad 1 \\ \hline 01100101 \end{array}$$

- Algorithm 2 – Start with the least significant bit, copy all of the bits up to and including the first 1 bit and then complementing the remaining bits.

$$\begin{array}{r} N = 01100101 \\ [N] = 10011011 \end{array} \quad \leftarrow \text{least significant bit}$$

Finite Number Representation

- Machines that use 2's complement arithmetic can represent integers in the range:

$$-2^{n-1} \leq N \leq 2^{n-1}-1$$

where n is the number of bits available for representing N .

- Note that $2^{n-1}-1 = (011\dots11)_2$ and

$$-2^{n-1} = (100\dots00)_2$$

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- For 2's complement more negative numbers than positive.
- For 1's complement two representations for zero.
- For a n bit number in base (radix) z there are z^n different unsigned values.

$$(0, 1, \dots, z^{n-1})$$

1's Complement Addition

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- Let's compute $(12)_{10} + (1)_{10}$.
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Step 1: Add binary numbers

Step 2: Add carry to low-order bit

$$\begin{array}{r} 01100 \\ + 00001 \\ \hline 01101 \\ 0 \\ \hline 01101 \end{array}$$

1's Complement Subtraction

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$$- (-1)_{10} = -(0001)_2 = 11110_2 \text{ in 1's complement}$$

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Step 1: Take 1's complement of 2nd operand.

$$\begin{array}{r} 01100 \\ - 00001 \\ \hline \end{array}$$

1's comp

$$\begin{array}{r} 01100 \\ + 11110 \\ \hline \end{array}$$

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$$\begin{array}{r}
 01100 \\
 - 11110 \\
 \hline
 \end{array}$$

Step 1: Take 1's complement of 2nd operand.

Add

$$\begin{array}{r}
 01100 \\
 + 11110 \\
 \hline
 \end{array}$$

Step 2: Add the binary numbers.

$$\begin{array}{r}
 101010
 \end{array}$$

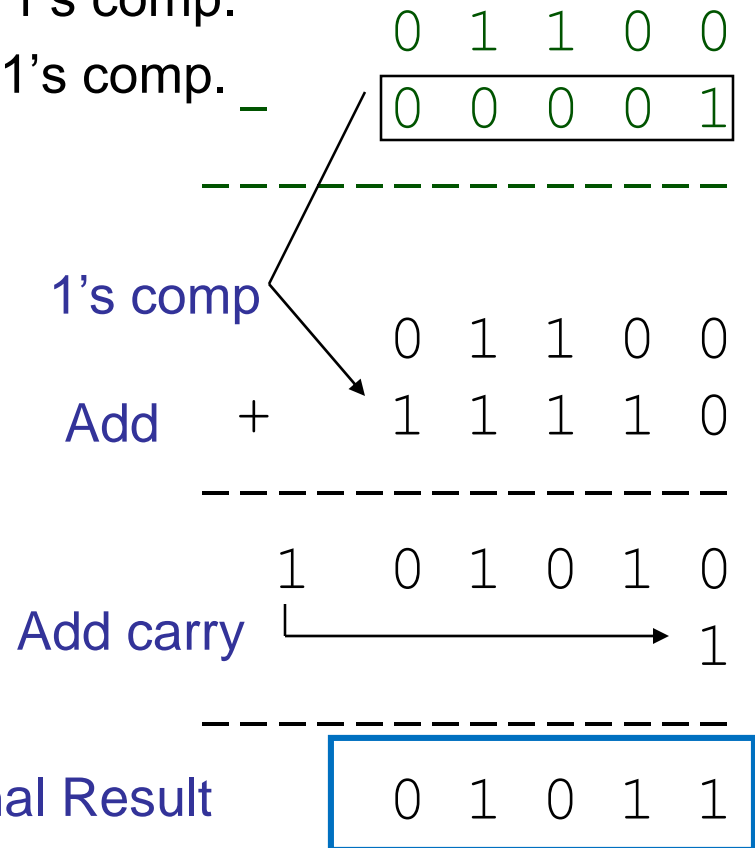
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2's Complement Addition

- Adding 2's complement numbers is easy.
- Let's compute $(12)_{10} + (1)_{10}$.
 - $(12)_{10} = +(1100)_2 = 01100_2$ in 2's comp.
 - $(1)_{10} = +(0001)_2 = 00001_2$ in 2's comp.

Step 1: Add binary numbers

Step 2: Ignore carry bit

$$\begin{array}{r} \text{Add} \quad + \quad \begin{array}{r} 01100 \\ 00001 \\ \hline \end{array} \\ \text{Final Result} \quad 0 \quad \boxed{01101} \\ \uparrow \\ \text{Ignore} \end{array}$$

2's Complement Subtraction

- Follow the 3 steps for subtraction.
- Let's compute $(12)_{10} - (1)_{10}$.
 - $(12)_{10} = +(1100)_2 = 01100_2$ in 2's comp.
 - $(-1)_{10} = -(0001)_2 = 11111_2$ in 2's comp.

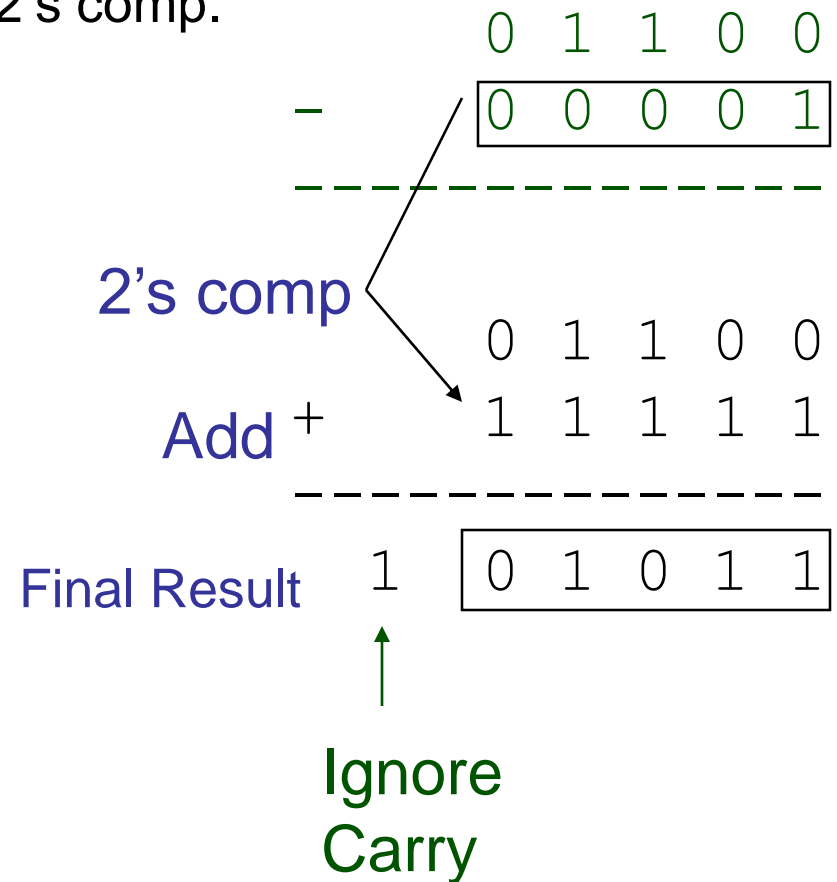
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Step 1: Take 2's complement of
2nd operand

Step 2: Add binary numbers

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2's Complement Subtraction: **Exercise 1**

Compute $(13)_{10} - (5)_{10}$ using the 2s complement form.

5 minutes
to complete this exercise!!

2's Complement Subtraction: Exercise 1

- Let's compute $(13)_{10} - (5)_{10}$.

$$\begin{aligned}(13)_{10} &= +(1101)_2 = (01101)_2 \\ (-5)_{10} &= -(0101)_2 = (11011)_2\end{aligned}$$

- Adding these two 5-bit codes...

carry

$$\begin{array}{r} \\ + \\ \hline 1 \end{array}$$

- Discarding the carry bit, the sign bit is seen to be zero, indicating a correct result. Indeed,

$$(01000)_2 = +(1000)_2 = +(8)_{10}$$

2's Complement Subtraction: Exercise 2

Compute $(5)_{10} - (12)_{10}$

5 minutes

to complete this exercise!!

Week 3 Lecture 4b

- Binary numbers arithmetic
- Signed binary numbers: signed magnitude, 1st complement and 2s complement
- 1s and 2s complement arithmetic
- Course web page:
https://ecs.wgtn.ac.nz/Courses/XMUT101_2021T1/
- kerese@ecs.vuw.ac.nz