EXAMINATIONS — 2005
END-YEAR

## COMP 202 <br> Formal Methods of Computer Science

Time Allowed: 3 Hours
Instructions: There are seven (7) questions, worth fourteen (14) marks each, making ninety-eight (98) marks in total.
Answer all the questions.
You may use printed foreign language dictionaries.
You may not use calculators or electronic dictionaries.

## Question 1.

(a) Describe the main idea behind the construction of a finite automaton from a nondeterministic finite automaton, such that both machines define the same language. [5 marks]
(b) Show that every finite language is regular.
(c) Let $M_{1}$ be the nondeterministic finite automaton with $\Lambda$ transitions defined by $M_{1}=\left(Q, \Sigma, \delta, q_{0}, F\right)$, where:

- $Q=\left\{S_{0}, S_{1}, S_{2}\right\}$
- $\Sigma=\{a, b\}$
- $q_{0}=S_{0}$
- $F=\left\{S_{2}\right\}$
and $\delta$ is given by the table:

| $\delta$ | $a$ | $b$ | $\Lambda$ |
| :--- | :--- | :--- | :--- |
| $S_{0}$ | $\left\{S_{0}\right\}$ | $\}$ | $\left\{S_{1}\right\}$ |
| $S_{1}$ | $\}$ | $\left\{S_{1}\right\}$ | $\left\{S_{2}\right\}$ |
| $S_{2}$ | $\left\{S_{2}\right\}$ | $\}$ | $\left\{S_{2}\right\}$ |

(i) Draw $M_{1}$.
(ii) Find a nondeterminstic finite automaton without $\Lambda$ transitions which accepts the same language as $M_{1}$.

## Question 2.

(a) Let $M_{1}$ and $M_{2}$ be nondeterministic finite automata with $\Lambda$ transitions, and let $L_{1}=\operatorname{Language}\left(M_{1}\right)$ and $L_{2}=$ Language $\left(M_{2}\right)$.

Draw nondeterministic finite automata with $\Lambda$ transitions which describe the following languages. Your diagrams should indicate how the states of your machines relate to the states of $M_{1}$ and $M_{2}$.
(i) $L_{1} \cup L_{2}$
(ii) $L_{1} L_{2}$
(iii) $L_{1} \cap \overline{L_{1}}$
(b) Let: $\Sigma=\{a, b\}$

$$
\begin{aligned}
& L_{1}=\operatorname{Language}\left(\mathbf{a}(\mathbf{a}+\mathbf{b})^{*}\right) \\
& L_{2}=\operatorname{Language}\left((\mathbf{b a})^{*}\right) \\
& L_{3}=\operatorname{Language}\left(\mathbf{a}^{*}+\mathbf{b}^{*}\right)
\end{aligned}
$$

Give a string which:
(i) is in $L_{1}$ but not in $L_{2}$ or $L_{3}$
(ii) is in $L_{2}$ but not in $L_{1}$ or $L_{3}$
(iii) is in $L_{3}$ but not in $L_{1}$ or $L_{2}$
(c) State Kleene's theorem.

## Question 3.

Consider the context free grammar

$$
\begin{array}{ll}
(1,2) & A \rightarrow a A \mid B \\
(3,4,5) & B \rightarrow a A b A|C| \Lambda \\
(6) & C \rightarrow b
\end{array}
$$

(a) Define $\Sigma, N$, and $S$ for this grammar. [2 marks]
(b) Draw a parse tree for the string $a a b b$.
(c) Show that the grammar is ambiguous.
(d) List the nullable nonterminals.
(e) List the unit productions.
(f) Find an equivalent grammar with no unit productions.

## Question 4.

Consider the context free grammar

$$
\begin{aligned}
(1,2) & S \rightarrow T S b \mid T \\
(3,4,5,6) & T \rightarrow a U a|b U b| a \mid b \\
(7,8) & U \rightarrow S \mid c T U
\end{aligned}
$$

(a) Explain what it means for a grammar to be in LL(1) form.
(b) Show that the above grammar is not in LL(1) form.
(c) Rewrite the grammar so that it is in LL(1) form.
(d) Complete the Parsel procedure whose heading appears below, for a recursive-descent parser to recognize the $U$ nonterminal. You may assume that procedures ParseS and ParseT for recognizing the other nonterminals are already written. The parameter ss, providing the input for the parser, is a sequence of terminal symbols. You do not need to return a parse tree.

```
procedure ParseU (in out ss)
```

begin
end

## Question 5.

(a) Define a pushdown automaton that accepts each of the following languages:
(i) The language consisting of strings over the alphabet $\{a, b\}$ with the same number of as and bs, but occurring in any order
(ii) The language $\left\{a^{m} b^{n} c^{m-n} \mid m \geq n>0\right\}$ [6 marks]
(b) Define a pushdown automaton that accepts the language generated by the following grammar, using a bottom-up (shift-reduce) strategy.

$$
\begin{aligned}
& S \rightarrow T U \mid \Lambda \\
& T \rightarrow a S a \mid S b \\
& U \rightarrow a U \mid a
\end{aligned}
$$

[4 marks]
(c) Briefly outline a proof that any language accepted by a pushdown automaton may be generated by a context free grammar.
[4 marks]

## Question 6.

(a) A Turing machine may be defined by a 6 -tuple ( $\left.Q, \Sigma, \Gamma, \delta, q_{0}, F\right)$.

Explain each of the following:
(i) $\Sigma$
(ii) $\Gamma$
(iii) $\delta$
(b) Let $L_{1}$ and $L_{2}$ be regular languages. Is it possible to write a program which decides whether $L_{1}$ and $L_{2}$ are the same language? Justify your answer.
(c)
(i) Define "universal Turing machine".
(ii) Briefly outline how a Turing machine may be encoded so as to provide the input for another Turing machine.
(iii) Discuss the relevance of the existence of universal Turing machines in proving that there are undecidable problems.

## Question 7.

You are given strings $s$ and $t$, both of length $n(|s|=|t|=n)$, and are required to determine whether $s$ is the reverse of $t\left(s=t^{R}\right)$.
(a) Give a formal specification (signature, precondition, and postcondition) for the problem.
(b) State three requirements for proving the correctness of a loop in a while-program using a loop invariant.
(c) The following fragment of program partly satisfies the specification: it is possible to determine from the final value of $i$ whether $s=t^{R}$.

```
begin
\(1 i:=n-1\);
\(2 j:=0\);
3 while \(i \geq 0\) and \(s[i]=t[j]\) do
\(4 \quad i:=i-1\);
\(5 \quad j:=j+1\)
6 end
```

An invariant that may be used to prove the loop correct is

$$
-1 \leq i<n \wedge j=n-i-1 \wedge(\forall k \in 0 . . j-1)(s[k]=t[n-k+1])
$$

(i) Show that the loop invariant holds at the start of the loop.
(ii) State an assertion that holds at the end of the loop, and show that it holds there.
(iii) Hence, give statements that can follow the loop (between lines 5 and 6) to complete the program, and show that your postcondition from part (a) is satisfied. [8 marks]

