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Student ID:
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# COMP261: Algorithms and Data Structures <br> <br> Term Test 3 

 <br> <br> Term Test 3}

9 May 2024

## Instructions

- Time allowed: 50 minutes
- Attempt all the questions. There are 50 marks in total.
- Write your answers in this test paper and hand in all sheets.
- If you think a question is unclear, ask for clarification.
- This test contributes $\mathbf{1 5 \%}$ of your final grade.
- You may write notes and working on this paper, but make sure your answers are clear.
- Only silent non-programmable calculators or silent programmable calculators with their memories cleared are permitted in this examination.
- No electronic dictionaries are allowed.
- Paper foreign to English language dictionaries are allowed.


## Sections

1. Network Flows
2. Centrality
3. Cycles and Spanning Trees

## Marks

TOTAL: $\square$

## SECTION A Network Flows

1. In a network flow problem, what constitutes the 'source'?
a) Vertex with the maximum capacity
b) Vertex with no outgoing edges
c) Vertex with no incoming edges
d) Vertex with the least capacity

2. How many constraints must a flow satisfy in a network flow problem?
a) One
b) Two
c) Three
d) four
$\square$
3. State True or False. In a standard network flow problem, under certain circumstances, a vertex v that is not a source or a sink can have a total in-flow that has a different value than its total outflow.
[1 mark]
$\square$
4. Consider the following graph


If the Ford-Fulkerson method identifies $\mathrm{S}->\mathrm{A}->\mathrm{C}->\mathrm{T}$ as the first augmentation path, are there any remaining augmentation paths within the residual graph? If so, list the vertices in a valid augmentation path in the box below along with the value of the corresponding flow. If there is no other augmentation path that can be found write 'none '.
5. Suppose that you are asked to implement the Edmond-Karp algorithm to find maximum flow through a network. Using the following pseudo-code of the algorithm answer the following questions.
[12 marks]

```
EdmondKarp(Graph, start, goal):
    // Build Residual Graph to include flow values and reverse edges
    for each edge e in \(G\) where \(e\) is of the form 〈fromVertex, toVertex, capacity〉
2 add e \(\leftarrow\langle\) fromVertex, toVertex, capacity, 0\(\rangle\) to the Residual graph \(/ /\) flow value \(=0\) initially
3 add the corresponding reverse edge \(\mathrm{e}^{\prime} \leftarrow\langle\) toVertex,fromVertex, 0,0\(\rangle\) to the Residual graph
maxflow \(\leftarrow 0\)
Repeat
        Use BFS to find a path \(P\) from start to goal in the Residual Graph such that
        all edges \(e\) in \(P\) have capacity (e)>0
7 if P exists
        //Find bottleneck:the capacity of the minimum capacity edge in \(P\)
8 pathFlow \(\leftarrow\) Bottleneck(P)
9 Output (P, pathFlow) as an augmentation path
10 maxflow \(\leftarrow\) maxflow + pathflow
11 Update the Residual graph to reflect the changes in the capacities of edges
            that constitue \(P\) and the corresponding reverse edges
13 else
14 output Maxflow
15 return
```

(a) [3 marks] In the following graph, if the first path discovered by BFS is $S->A->B->T$, which edges in the residual graph will be updated at step 11 of the pseudo-code? Also list the updated capacity and flow for each updated edge.

(b) [8 marks] For the path $S->A->B->T$ identified above, trace out what the BFS will do at each iteration in terms of:

- contents of the Queue
- contents of the Backpointers Map

Note: If at a given iteration more than one node is to be added to the queue, add them in alphabetical order. The first iteration has been done for you.

Queue: S

## Iteration 1:

Queue: A C

Backpointers: $\langle\mathrm{A} \rightarrow$ Edge S-to-A $\rangle,\langle\mathrm{C} \rightarrow$ Edge S-to-C $\rangle$

## Iteration 2:

Queue:
Backpointers:

## Iteration 3:

Queue:
Backpointers:

Iteration 4:

Queue:

Backpointers:
(c) [1 mark] When running Ford-Fulkerson as per the implementation, how does the algorithm know when to terminate?
6. Name the constraint that the flow identified in figure 2, for the network depicted in figure 1, fails to satisfy.


Figure 1


Figure 2

## SECTION B Centrality

7. Trace out the steps of the PageRank algorithm for the given web graph using the provided pseudo-code with the parameter iter set to 1 and dampingfactor set to 85 .
```
ComputePageRank(Graph, iter, dampingfactor):
    nNodes \(\leftarrow\) get count of nodes in the graph
    for each node n in Graph
        set pageRank \((\mathrm{n}) \leftarrow 1.0 / \mathrm{nNodes}\)
    count \(\leftarrow 1\)
    Repeat
        randomjumpShare \(\leftarrow\) (1-dampingfactor)/nNodes
        sinkShare \(\leftarrow 0\)
        for each node n in the Graph that is a sink node (has no outlinks)
        sinkShare \(\leftarrow\) sinkShare + dampingfactor \(\times\) (pageRank( n )/nNodes)
        for each node n in the Graph
        neighbourShare \(\leftarrow 0\)
        for each neighbour \(b\) of \(n\)
            neighbourShare \(\leftarrow\) neighbourShare + pageRank(b)/outdegree(b)
            newPageRank \((\mathrm{n}) \leftarrow\) sinkShare + randomjumpShare \(+\mathrm{d} \times\) neighbourShare
        Update pageRank with newpageRank
        count++
```

    Until count > iter
    
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$\square$
8. Which of the following will you use to assess the vulnerability of critical infrastructure networks, such as water distribution systems or power grids, by identifying nodes that, if disrupted, could lead to cascading failures?
[1 mark]
a) Degree centrality
b) Closeness centrality
c) Betweenness centrality
d) Network flows

9. What kind of centrality would you want to analyze in a graph if you wanted to inject information that flows through the shortest path in a network and have it spread quickly?
a) Degree
b) Group
c) Closeness
d) Betweenness

10. Select all of the following problems where closeness would be a valuable centrality measure.
[2 marks]
a) Optimize Bus stops locations
b) Optimize Bus routes
c) Optimize Traffic signal timings
d) Optimize Bus stop timings along a route
$\square$

## SECTION C Cycles and Spanning Trees

11. When using DFS for cycle detection, what does the presence of a back edge indicate? [ $\mathbf{1}$ mark]
a) The graph is acyclic
b) A cycle is present in the graph
c) The graph is connected
d) The graph is disconnected
$\square$
12. Which of the following data structures is typically used in Prim's algorithm to efficiently select the next edge to add to the spanning tree?
a) Priority queue
b) Stack
c) Queue
d) Array

13. In Kruskal's algorithm, which of the following operations is performed to efficiently determine whether adding an edge creates a cycle in the spanning tree?
[1 mark]
a) Breadth-first search
b) Depth-first search
c) Find operation with disjoint data structure
d) Dijkstra's algorithm
$\square$
14. If a graph has $n$ vertices, what is the maximum number of edges it can have without forming a cycle?
a) $\mathrm{n}-1$
b) $n$
c) $\mathrm{n}-2$
d) $2 n$

15. State True or False: If a graph has multiple edges with the same weight, Prim's algorithm may produce different minimum spanning trees depending on the order in which edges are processed.
[1 mark]
$\square$
16. Using Prim's algorithm find the cost of the minimum cost spanning tree of the following graph. You should show the edges and their weights (e.g., AC 6) in the order of being added into the tree by the algorithm considering A as the starting node.
[6 marks]
Note: The cost of a spanning tree equals the sum of the costs of all its edges.



Using the provided pseudo-code, answer the following questions related to the call to CycleDetection(g):
a) List the sequence of the calls made to isCyclicDFS() starting with the initial call isCyclicDFS (g,

A, visited, -1).
b) Identify the call when the algorithm detects the presence of a cycle. Provide an explanation for your answer.
c) For each call also identify which node will be marked as visited.

Note: If at a given iteration a node has more than one neighbours, traverse them in alphabatic order.

```
function CycleDetection(Graph graph):
    nNodes }\leftarrow\mathrm{ get count of nodes in the graph
    boolean visited [nNodes] // all initialized to false as none of the nodes have been visited
    for each node v in the graph
        if (! visited [v])
            if (isCyclicDFS(graph, v, visited, , -1))
                return true
    return false
end function
function isCyclicDFS(Graph, start, visited [], parent)
    set visited [start] to true
    for each neighbour n of node start
        if (! visited [n])
            if (isCyclicDFS(graph, n, visited, start ))
            return true
        else if (n != parent)
            return true
    return false
end function
```

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18. Consider the following forest of trees maintained as a disjoint-set data structure.

Tree 1:

Tree 2:


Tree 3:


Using the pseudo-code below, present the forest graphs obtained after each of the following two merge operations performed in a sequence: Union ( $B, E$ ) followed by Union (I, D).

```
Union(x, y ):
    xroot }\leftarrow\operatorname{Find}(x
    yroot }\leftarrow\mathrm{ Find(y)
    if (xroot == yroot)
        return
    else
        if (xroot.depth < yroot.depth)
            xroot.parent }\leftarrow\mathrm{ yroot
            remove xroot from forest
        else
            yroot.parent \leftarrowxroot
            remove yroot from forest
            if (xroot.depth == yroot.depth)
                xroot.depth ++
```


## SPARE PAGE FOR EXTRA ANSWERS

Cross out rough working that you do not want marked.
Specify the question number for work that you do want marked.

