#### **STUDENT:**

## **QUESTION 1**

Part	Description	Marks	Comment
(a) 2 marks	<ul> <li>No electromagnetic interference.</li> <li>Therefore, no need for error detection/correction.</li> </ul>		
(b) 2 marks	<ul> <li>Error detection code added to the frame before sending.</li> <li>Code recomputed and compared with received EDC.</li> </ul>		
(c) 2 marks	<ul> <li>Record address as associated with interface.</li> <li>Broadcast on all links except on the one received.</li> </ul>		

Part	Description	Marks	Comment
(a) 2	4 bytes.		
marks	$\square$ 16 bytes.		
(b) 4	$\square$ (i) = C		
marks	□ (i) MSB is 110.		
	$\square$ (ii) = A		
	$\square$ (ii) MSB is 0.		
(c)(i) 2	□ 32-26=8-2=6		
marks	$\Box$ 2^6=64.		
	Said STATE so working does not need to be shown.		
(c)(ii) 2	$\square$ Fourth subnet (count from 0 so 3), 11000000		
marks	□ 129.254.128.192		
	Note that $1^{\text{st}}$ subnet is 00000000 so 129.254.128.0, $2^{\text{nd}}$ subnet is		
	01000000 so 129.254.128.64, 3 <sup>rd</sup> subnet is 10000000 so		
	129.254.128.128.		
	Said STATE so working does not need to be shown.		
(c)(iii) 2	$\square$ Fourth subnet (count from 0 so 3), 11111111		
marks	□ 129.254.128.255		
	Broadcast addresses for other subnets are: 129.254.128.		
	Said STATE so working does not need to be shown.		

### STUDENT:

Part	Description	Marks	Comment
(a) 4	□ GET		
marks	□ Host		
	□ Operation		
	□ Parameter		
(b) 4	D POST		
marks	$\Box$ Service (/deposit).		
	□ Account as XML		
	□ Amount as XML		
(c) 4	D POST.		
marks	□ /withdraw		
	□ Account as XML		
	□ Amount as XML		
(d) 4	□ With POX, the service always returns 2000K irrespective of		
marks	the error.		
	□ The programmer has to parse the response to extract error		
	information.		
	□ With CRUD services. The HTTP status codes are used to		
	indicate the type of error.		
	□ The response contains additional information about the error.		
(e) 4	□ Timestamp is cheaper to compute than a hash.		
marks	□ Hash requires doing a computation over entire reply.		
	Granularity of changes that can be monitored.		
	□ May change more rapidly than the resolution of the		
	timestamp.		

Part	Description	Marks	Comment
(a) 2 marks	<ul> <li>Something like http://www.bank.com/account.</li> <li>Minimises coupling with clients by hiding server implementation details (encapsulation)</li> </ul>		
(b) 2 marks	<ul> <li>Allows different clients to be used,</li> <li>For example a browser rather than a program might be used to interact with the service so returning XHTML might be a better choice.</li> </ul>		
(c) 5 marks	<ul> <li>Allows programmers to take advantage of standard rendering libraries.</li> <li>Thereby reducing the work of building clients and servers.</li> </ul>		
(d) 6 marks	<ul> <li>Each state is listed.</li> <li>Logout leads to no state.</li> <li>Customer is able to getBalance, deposit, withdraw, logout for the rest for at least one example.</li> <li>Appropriate links for the remaining states.</li> <li>Attempted to use LINK type to indicate flow.</li> </ul>		
(e) 3 marks	<ul> <li>Login provides getBalance and logout links.</li> <li>GetBalance provides withdraw and deposit and logout.</li> <li>Withdraw and deposit provide GetBalance and logout.</li> </ul>		
(f) 5 marks	<ul> <li>When links are returned to clients that have not been updated, the clients just ignore them.</li> <li>This means that old clients can keep running without needing to be updated.</li> <li>This improves the robustness of the application by removing potential fragilty.</li> <li>That is in the previous approaches the client would need to be updated to work with the application each time that a protocol is changed.</li> <li>Additionally the semantic information returned with responses allow clients to be flexible in how they render representations.</li> </ul>		

Part	Description	Marks	Comment
(a) 4 marks	<ul> <li>(1) Client retrieves representation from a nearby cache without having to traverse the whole network to the origin server thereby reducing bandwidth usage;</li> <li>(2) Client retrieves representation from a proxy that is closer in terms of latency thereby improving performance;</li> <li>(3) Instead of all requests going to a single server, the load is spread out over multiple proxy servers.</li> <li>(4) Where it is a transient failure, the proxy can respond to requests from clients even if the origin server is unavailable.</li> </ul>		
(b) 6 marks	<ul> <li>Local cache.</li> <li>Stores representations from many origin servers on behalf of a single user agent, application or machine.</li> <li>Proxy cache.</li> <li>Proxy cache stores representations from many origin servers on behalf of many consumers. Can be hosted both inside the corporate firewall and outside.</li> <li>Reverse proxy.</li> <li>Stores representations from one origin server on behalf of many consumers. Reverse proxies are located in front of an application or web server.</li> </ul>		
(c) 2 marks	<ul> <li>Requires server to maintain knowledge of which clients it is in conversation.</li> <li>Allows it to inform them of changes to its state.</li> </ul>		
(d) 3 marks	<ul> <li>Only the client's local cache are allowed to store the data.</li> <li>Can be stored for up to an hour.</li> <li>Would be used where the response is somehow private to the consumer and so should not be stored in a public cache and can be up to an hour out of date (caching allows better performance).</li> </ul>		
(e) 5 marks	<ul> <li>Fresh if it has not been stored in the cache longer than a specified lifetime.</li> <li>Has nothing to do whether it has diverged from the original or not.</li> <li>The original may have changed but the cached representation might still be fresh.</li> <li>A stale cached representation is one that has exceeded its lifetime.</li> <li>This might have diverged but equally the representation at the origin server might be just the same.</li> <li>You can only find this out by revalidating the cached representation.</li> </ul>		

Part	Description	Marks	Comment
(a) 2 marks	<ul> <li>An event (as specified in the book) is a significant change in the state of the resource at any particular point in time,.</li> <li>In this case it would be a change in temperature.</li> </ul>		
(b) 3 marks	<ul> <li>Tight requires client and server to be present at the same instant of time in order for communication to take place.</li> <li>Loose doesn't require both to be present at the same time for communication to take place.</li> <li>Loose improves reliability of the application.</li> </ul>		
(c) 5 marks	<ul> <li>Publisher is polled by the consumer.</li> <li>Consumer checks at regular intervals if there are any new events.</li> <li>The shorter the polling period the more up to date the view of the events.</li> <li>But the greater bandwidth used.</li> <li>Worse performance of the client as it has to do more work.</li> </ul>		
(d) 5 marks	<ul> <li>Feed is fetched as a single document therefore do want it to grow forever.</li> <li>Otherwise everytime we want to poll we will incur large bandwidth costs.</li> <li>Therefore, break the feeds up into chunks that are smaller files and less costly to download.</li> <li>One approach is simply archive at a regular interval but this does not place an upper bound on the size of a feed.</li> <li>Another approach is to simply set a maximum size or number of entries and archive when it reaches this size.</li> </ul>		
(e) 5 marks	<ul> <li>Client finds current event feed. Reads events to see if any that it has seen before.</li> <li>No? got to the previous archive and repeat the process until eventually a previously processed event is found.</li> <li>When found, process all events in the feed.</li> <li>Then follow the next-archive link to the next oldest archive and process these events. Continue until all events are processed.</li> <li>This should be expensive because reading archives twice.</li> <li>However, if the feeds are cached it is likely that they will be stored in a local cache which does not require any requests to pass over the network to the origin server.</li> </ul>		

Part	Description	Marks	Comment
(a) 4 marks	<ul> <li>Need to construct a function of degree 2.</li> <li>For example, f(x) = 42 + 3x.</li> <li>Hand out shares (= number of people to share secret with), for example (1, f(1)), (2, f(2), (3, f(3)), (4, f(4)).</li> <li>Need two people in order to work out what the secret was through interpolation, one is not sufficient.</li> </ul>		
(b) 3 marks	<ul> <li>Malice sends Bob "Hi, I'm Alice" and gives Bob Malice's public key.</li> <li>Malice intercepts Bob's challenge.</li> <li>Malice sends Bob N encrypted using Malice's private key.</li> </ul>		
(c) 2 marks	<ul> <li>Authentication is about proving who you are whereas.</li> <li>Authorisation is about proving your right to do something.</li> </ul>		
(d) 3 marks	<ul> <li>Ticket for service issued to the client has a time to live.</li> <li>An attacker could extend the time to live by making the application server's clock run slow.</li> <li>This would allow them to capture the ticket issued to the client and replay it a many times as it would like.</li> </ul>		
(e) 2 marks	<ul> <li>First, the service doesn't have to manage its own authentication system.</li> <li>Second, it provides a single sign on service freeing the user from having to remember a large number of usernames and passwords.</li> </ul>		
(f) 3 marks	<ul> <li>The reason isn't confidentiality,</li> <li>The main reason is that a malicious service might redirect a user to a fake version of their OpenID provider</li> <li>In order to steal user credentials.</li> </ul>		
(g) 5 marks	<ul> <li>Server – service hosting the protected resources – this would be the Bank. Issues temporary credential for access.</li> <li>Resource owner – the customer in this case who has control over the resource.</li> <li>Authenticates itself and chooses whether credentials will be issued to the client.</li> <li>Client – the service needing access to the protected resource.</li> <li>Must request authorisation from the Bank for access to resources and bank redirects it to the resource owner.</li> </ul>		