

EXAMINATIONS — 2013  
END-OF-YEAR

**SWEN 224**  
**Formal Foundations**  
**of Programming**

**Time Allowed:** 3 Hours

- Instructions:**
- Answer all **four** questions.
  - The exam will be marked out of one hundred and eighty (180).
  - Calculators ARE NOT ALLOWED.
  - Non-electronic foreign language dictionaries are allowed.
  - No other reference material is allowed.

## Question 1. Assertions

[45 marks]

Let  $A$  be an array of integers indexed from 1 to  $|A|$  and  $B$  be an array of integers indexed from 1 to  $|B|$ . Let  $x, y$  and  $z$  be integers.

(a) Consider the following assertion:

$$x < y \wedge z = y$$

(i) [1 mark] State in words what the assertion means.

For each of the following, say whether the assertion is true or false:

(ii) [1 mark]  $x = 2, y = 3, z = 3$

(iii) [1 mark]  $x = 3, y = 3, z = 3$

(b) Consider the following assertion:

$$\forall j \bullet 1 \leq j \leq |A| \rightarrow A[j] > 50$$

(i) [2 marks] State in words what the assertion means.

For each of the following arrays, say whether the assertion is true or false:

(ii) [1 mark]  $A = ()$

(iii) [1 mark]  $A = (101, 102, 49)$

(c) Consider the following assertion

$$\left( \forall j \bullet 1 \leq j < |A| \rightarrow A[j] \leq A[j + 1] \right) \wedge \left( \sum_{i=1}^{|A|} A[i] \leq 0 \right)$$

(i) [4 marks] State in words what the assertion means.

For each of the following arrays, say whether the assertion is true or false.

(ii) [1 mark]  $A = (-10, 0, 1, 3, 3)$

(iii) [1 mark]  $A = (-10, 9, 0)$

(iv) [1 mark]  $A = ()$

(v) [1 mark]  $A = (-1, 0)$

**(d) Write an assertion formalising each of the following statements:**

**(i)** [4 marks]  $x$  is greater than  $y$  and less than  $z$ .

**(ii)** [4 marks]  $A$  contains at least one occurrence of  $d$ .

**(iii)** [4 marks] If  $A$  contains  $z$  then it does not contain  $y$ .

**(iv)** [4 marks] The sum of the elements in  $A$  is greater than the largest element in  $B$ .

**(v)** [4 marks] No element in  $A$  is in  $B$ .

**(e) Program contracts.**

**(i)** [5 marks] Write a formal contract, a pre-condition and post-condition, for the following program.

The program takes two arrays,  $A$  and  $B$ , containing the lengths and widths respectively of a sequence of rectangles. So  $A[i]$  and  $B[i]$  are the length and width of rectangle  $i$  (the  $i$ th rectangle). The program outputs an array,  $C$ , containing the area of the rectangles. Note the area of a rectangle is defined by  $area = length \times width$ .

**(ii)** [5 marks] Write a formal contract, a pre-condition and post-condition, for the following program.

The program inputs,  $P$ , an array of pairs. When  $i$  is an index of  $P$  then the first element of pair  $P[i]$  is written as  $first(P[i])$  and is a student name. The second element of  $P[i]$  is written as  $second(P[i])$  and is the integer grade for the student. The program outputs:

1.  $SB$  - the set of students who have grade "B", that is their numeric grade is greater than or equal to 70 and less than 75.
2.  $CB$  - the number of students that have grade "B".

## Question 2. Verifying Programs

[45 marks]

Each question has a program and a post-condition. Use Hoare Logic to compute the missing conditions, simplifying your answers where appropriate.

Variables  $x, y, z, w \dots$  are integers and  $A$  is an array of integers.

(a) [6 marks] What are the conditions: **P**, **P1** and **P2**?

<b>P</b>	...
	$y := 3x;$
<b>P1</b>	...
	$y := y + 2x;$
<b>P2</b>	...
	$x := y + 2;$
<b>Q</b>	$\{x \times y > 20\}$

(b) [6 marks] What are the conditions: **P**, **P1**, **P2** and **P3**?

<b>P</b>	...																
<table border="1"><tr><td><b>P1</b></td><td>...</td></tr><tr><td></td><td><math>x := x - 1;</math></td></tr><tr><td><b>Q</b></td><td><math>\{10 &gt; x + y\}</math></td></tr></table>	<b>P1</b>	...		$x := x - 1;$	<b>Q</b>	$\{10 > x + y\}$	<table border="1"><tr><td><b>P2</b></td><td>...</td></tr><tr><td></td><td><math>y := -y;</math></td></tr><tr><td><b>P3</b></td><td>...</td></tr><tr><td></td><td><math>x := 12</math></td></tr><tr><td><b>Q</b></td><td><math>\{10 &gt; x + y\}</math></td></tr></table>	<b>P2</b>	...		$y := -y;$	<b>P3</b>	...		$x := 12$	<b>Q</b>	$\{10 > x + y\}$
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<b>If</b> $(0 \leq x \wedge 0 \leq y)$ <b>then</b> $x := x - 1;$ <b>else</b> $y := -y; x := 12;$ <b>fi</b>																	
<b>Q</b>	$\{10 > x + y\}$																



(e) [3 marks] For a while program with guard **B** and loop body **S**:

```
while B do { S }
```

state the three loop conditions using pre-condition **P** and post-condition **Q**.

(f) Consider the following program:

```
int i := 0;
int x := 0;
while (i < y) then {
  x := x + b;
  i := i + 1;
}
```

(i) [3 marks] Given the loop pre-condition  $\mathbf{P} = \{i = 0 \wedge x = 0 \wedge 0 < y \wedge b \times y \leq 16\}$ , loop post-condition  $\mathbf{Q} = \{x \leq 16\}$  and the program above define:

(I) the loop invariant **L**

(II) the variant **V** (a value that decreases and is bounded below)

that will satisfy the three loop conditions.

(ii) [3 marks] Prove your loop invariant satisfies the three loop conditions for the given program, pre-condition and post-condition.

(g) The program below takes two arrays *A* and *B* of the same size and computes *s* the difference between the sum of the elements in *A* and the sum of the elements in *B*.

```
int i := 0;
int s := 0;
while i < |A| then {
  i := i + 1;
  s := s + A[i] - B[i];
}
```

(i) [6 marks] Define the following conditions:

(I) the post-condition to the while loop, **Q**

(II) the pre-condition to the while loop, **P**

(III) the loop invariant **L**

(ii) [6 marks] Prove your loop invariant, pre-condition and post-condition satisfy the three loop conditions for the given program.

### Question 3. Regular Languages

[45 marks]

(a) [8 marks] Write a *Regular Expression* to describe each of the following languages:

- (i) All strings over  $\{a, b\}$  which start and end with an  $a$  and do not contain two consecutive  $b$ 's.
- (ii) All strings over  $\{a, b, c\}$  containing exactly one  $a$  or exactly one  $b$  (or both). For example, this language includes the strings  $a$ ,  $cbaccaaa$ ,  $cbbba$  and  $ccaccbc$ , but not  $aacc$ ,  $ccbbccc$  or  $aabbb$ .

(b) [6 marks] Consider the language consisting of all strings over  $\{a, b, c\}$  in which no  $a$  is followed immediately by a  $b$ , no  $b$  is followed immediately by a  $c$  and no  $c$  is followed immediately by an  $a$  (i.e. no string contains  $ab$ ,  $bc$  or  $ca$  as a substring).

- (i) Draw a transition diagram for a DFA that accepts this language.
- (ii) Give a formal definition for your DFA in part (i).

(c) [6 marks] Draw a transition diagram for an  $NFA_\epsilon$  that accepts the language defined by the regular expression  $a(ab|b^*(a|c))^*a$ .

(d) Consider the NFA  $M = (Q, q_I, A, N, F)$ , where:

$$\begin{aligned} Q &= \{1, 2, 3, 4, 5, 6\} \\ q_I &= 1 \\ A &= \{a, b, c\} \\ N &= \{(1, a, 1), (1, a, 2), (1, a, 3), (2, a, 2), (2, a, 4), (3, a, 3), (3, b, 3), (3, b, 5), \\ &\quad (4, a, 4), (4, b, 4), (4, c, 6), (5, a, 5), (5, b, 5), (5, c, 5), (5, b, 6)\} \\ F &= \{5, 6\} \end{aligned}$$

- (i) [4 marks] Draw a transition diagram for  $M$ .
- (ii) [5 marks] Show the sequence of configurations that  $M$  passes through in accepting the input  $aababc$ .  
Note that you should show *all* states that  $M$  may be in after accepting part of the input.
- (iii) [8 marks] Draw a transition diagram for the DFA obtained by applying the *subset construction* to  $M$ . Show the correspondence between states of the DFA and those of the NFA. You only need to show reachable states.

(e) [8 marks] Given a DFA  $M = (Q, q_I, A, N, F)$  which accepts a language  $L$ , show how to construct a DFA,  $M' = (Q', q'_I, A', N', F')$ , which accepts the language  $L'$  consisting of all *prefixes* of strings in  $L$ , i.e.  $L' = \{\alpha \in A^* \mid \exists \beta \bullet \alpha\beta \in L\}$ .

Give a brief argument to show that the resulting DFA does in fact accept the required language.

## Question 4. Context-Free Languages

[45 marks]

(a) Suppose you want to use a simple language to describe the arrangement of shapes on a screen, where shapes can be nested within other shapes. Each shape is either a square or a circle, and is described by the position of its centre (a pair of natural numbers) and, optionally, a list of shapes contained within it. For example, the first example below describes a circle containing a square and another circle; the second example describes a square containing the circle from the first example, along with another circle:

- circle 80 60 [square 70 50, circle 90 90]
- square 50 50 [circle 80 60 [square 70 50, circle 90 90], circle 20 20]

(i) [10 marks] Write a *Context Free Grammar* to describe this language.

You are not required to give a full formal definition of this grammar — just write the list of rules. In writing this grammar, you should treat number as a terminal symbol and not define it further.

(ii) [5 marks] Draw a parse tree for the first example above using your grammar and assuming that 80, 60, etc. are instances of number.

(iii) [10 marks] Is your grammar LL(1)? If not, turn it into one that is.

Prove that the resulting grammar is LL(1).

(b) The following is part of a grammar for simple set expressions, where  $C$  stands for set comprehension and is not specified here:

$$E \rightarrow C \mid E \cup E \mid E \cap E \mid E - E \mid ( E )$$

(i) [6 marks] Explain, using an example, why this grammar is ambiguous.

(ii) [8 marks] Write an equivalent, unambiguous grammar, treating  $\cup$ ,  $\cap$  and  $-$  as *left-associative*, and giving  $\cup$  and  $-$  the same precedence, and  $\cap$  higher precedence than  $\cup$  and  $-$ .

(c) [6 marks] Let  $G = (V_N, V_T, S, P)$  be a Context Free Grammar which defines a language  $L$ .

Show how you can extend  $G$  to obtain a grammar  $G'$  which defines a language  $L' = \{(\alpha) \mid \alpha \in L\}$ , i.e. the language consisting of strings from  $L$  enclosed in a pair of brackets.

Give a brief argument to show that your grammar defines exactly the required language.

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