

SWEN430 - Compiler Engineering

Lecture 13 - Bytecode Generation II

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Generating Invoke Bytecodes

`invokevirtual` [1 byte op][2 bytes index]

Invoke method on a receiver of class type. The method and receiver types are located in the constant pool at the given index.

`invokeinterface` [1 byte op][2 bytes index]

Invoke method on a receiver of interface type. The method and receiver types are located in the constant pool at the given index.

`invokestatic` [1 byte op][2 bytes index]

Invoke static method. The method and receiver types are located in the constant pool at the given index.

`invokespecial` [1 byte op][2 bytes index]

Invoke special method (e.g. constructor). The method and receiver types are located in the constant pool at the given index.

Generating Invoke Bytecodes (Cont'd)

```
class Test {  
    Test(int x) { }  
    int f(String s, int i) {  
        return 1;  
    }  
  
    static void m(String[] s) {  
        Test t = new Test(123);  
        t.f(s[0],2);  
    }  
}
```

```
static void m(String[] s):
```

```
Code:
```

```
0:   new Test  
3:   dup  
4:   bipush 123  
6:   invokespecial Test.<init>:(I)V  
9:   astore_1  
10:  aload_1  
11:  aload_0  
12:  iconst_0  
13:  aaload  
14:  iconst_2  
15:  invokevirtual Test.f:(L...;I)I  
18:  pop  
19:  return
```

- Receiver pushed on stack first (line 10)
- Parameters pushed on stack next in order (lines 13-14)
- Return value is popped afterwards since its not used (line 18)

Determining Maximum Stack Height

- Must determine **maximum stack height** of each method
- Java Compiler can calculate the **stack difference** for each bytecode:
- Examples:

Bytecode	Stack Difference	Bytecode	Stack Difference
bipush	+1	pop	-1
iload X	+1	lload X	+2
iadd	-1	dadd	-2
iaload	-1	daload	0
ineg	0	d2f	-1
invokevirtual	???		

- Then, it traverses bytecode sequence determining the max height
- How should we deal with branching?

Determining Maximum Stack Height (cont'd)

```
int f(java.lang.String[]);
```

```
Code:                # diff # height
0:   aload_1          #  +1  #  1  #
1:   ifnull  12        #  -1  #  0  #
4:   getstatic System.out #  +1  #  1  #
7:   ldc  "Hello_World" #  +1  #  2  #
9:   invokevirtual println:(Ljava/lang/String;)V
      #  -2  #  0  #
12:  iconst_0          #  +1  #  1  #
13:  istore_2           #  -1  #  0  #
14:  iload_2            #  +1  #  1  #
15:  aload_1            #  +1  #  2  #
16:  arraylength        #   0  #  2  #
17:  iadd                #  -1  #  1  #
18:  istore_2           #  -1  #  0  #
19:  iload_2            #  +1  #  1  #
20:  ireturn            #  -1  #  0  #
```

- Hence, the maximum stack height for this method is 2

Line Number Information

```
int f(int x, int y) {           int f(int, int);
    int r = x / y;             0:   iload_1
    return r;                  1:   iload_2
}                               2:   idiv
                               3:   istore_3
                               4:   iload_3
                               5:   ireturn
                              LineNumberTable:
                               line 3: 0
                               line 4: 4
```

```
Exception in thread "main" java.lang.ArithmeticException: / by zero
    at Test.f(Test.java:3)
    at Test.main(Test.java:8)
```

- Can associate **line number information** with Bytecode
- `LineNumberTable` attribute is for this (see JVM Spec §4.7.8)
- Then can see **where** exceptions occur in original source file

Type Information

```
int f(int, int);  
  0:   iload_1  
  1:   iload_2  
  2:   idiv  
  3:   istore_3  
  4:   iload_3  
  5:   ireturn  
  ...
```

```
int f(int x, int y) {  
    int r = x / y;  
    return r;  
}
```

LocalVariableTable:

Start	Length	Slot	Name	Signature
0	6	0	this	Ltest;
0	6	1	x	I
0	6	2	y	I
4	2	3	r	I

- LocalVariableTable identifies **type** of each slot at each point

Exception Handlers

```
String f(Integer i) { String f(Integer);  
  try { Code:  
    return i.toString(); Stack=1, Locals=3  
  } catch (Exception e) { 0: aload_1  
    return ""; 1: invokevirtual Integer.toString()  
  } 4: areturn  
  5: astore_2  
  6: ldc ""  
  8: areturn  
  Exception table:  
    from to target type  
      0 4 5 Class Exception
```

- Exception handlers implemented as table rows:
 - » Range of bytecodes, destination and class
 - » Range cannot include handler itself

Exception Handlers Table

- Java Compiler generates exception handlers such that:
 - » Either exception handler ranges are **disjoint**, or one is a **subrange** of the other
 - » Exception handler code is never **within its own range**
 - » Entry for exception handler only via exception (not via e.g. `goto`)
- Surprisingly, these restrictions **not enforced** by bytecode verifier
 - » Because not considered a threat to integrity of JVM
 - » Still require that e.g. every nonexceptional path to handler has a single object on the operand stack, etc
 - » See JVM Specification, §3.10

Another Example

```
String f(Integer i) {  
    try {  
        return i.toString();  
    } catch (NullPointerException e) {  
        return "null";  
    } catch (Exception e) {  
        return "";  
    }  
}}  
  
String f(Integer);  
Code:  
Stack=1, Locals=3, Args_size=2  
0:   aload_1  
1:   invokevirtual Integer.toString:()  
4:   areturn  
5:   astore_2  
6:   ldc   "null"  
8:   areturn  
9:   astore_2  
10:  ldc   ""  
12:  areturn  
Exception table:  
from  to  target type  
    0    4    5    Class NullPointerException  
  
    0    4    9    Class Exception
```

- Multiple Exception handlers are triggered **in order of appearance**

Translating While: Type Representation

WHILE	JVM	Notes
<code>bool</code>	<code>JvmTypes.Bool</code>	Compiles to <code>int</code> in some cases.
<code>int</code>	<code>JvmTypes.Int</code>	
<code>T[]</code>	<code>List<Object></code>	Boxing / unboxing
<code>{T f, ... }</code>	<code>Map<String, Object></code>	Boxing / unboxing

- **Suggested mapping** from WHILE to JVM types (above)
- You will encounter an **impedance mismatch** though!
- For example, how to **translate** “`x = [1, 2, 3]`” ... ?

Translating While: Boxing / Unboxing

WHILE	JAVA
<pre>type Rec is {int g} int y = 1; Rec xs = { f: y+1 };</pre>	<pre>int y = 1; Map<String, Object> xs = ...; xs.put("f", y+1);</pre>

- **Compiles** in JAVA ...

... but on JVM is more **complicated** ...

Why?

Translating While: Boxing / Unboxing

```
0:  iconst_1
1:  istore_1           // y = 1
2:  new    java/util/HashMap
5:  dup
6:  invokespecial  java/util/HashMap."<init>": ()V
9:  astore_2
10: aload_2
11: ldc    "f"
13: iload_1
14: iconst_1
15: iadd           // y + 1
16: invokestatic  java/lang/Integer.valueOf: (I);
19: invokeinterface java/util/Map.put: (Object;Object;);
24: pop
25: return
```

Translating While: Cloning

WHILE	JAVA
<pre>int [] xs = [1, 2]; int [] ys = xs; ys[0] = 3;</pre>	<pre>List xs = new ArrayList(); xs.add(1); xs.add(2); List ys = xs; ys.set(0, 3);</pre>

- Above is **broken** ... *why?*
- **Remember:** WHILE has **value semantics** for arrays ...

... but JVM does **not!**