

Shuffle a list

Given a list, put items into a random order

23,22,49,25,43,23,5,31,43,27,21,45,43,16,5,21,18,27,39,18,21,7,42,28,21,19

- For each position, grab a random item and put it in that position
 - `add(position, remove(random))`
- vs
- `swap [set(position, set(index, get(position))]` or `Collections.swap(...)`

- Use the built-in shuffle!
 - `Collections.shuffle(list)`

Shuffle a list

n times {

- For each position from $n-1$ to 0 ,
 - choose a random index \leq position $n \times O(1)$
 - item = remove(index) $n \times O(n)$
 - add(position, item) $n \times O(n)$

Total: $O(n^2)$

n times {

- For each position from $n-1$ to 0 ,
 - choose a random index \leq position $n \times O(1)$
 - swap(index, position) $n \times O(1)$

Total: $O(n)$

Combinations

Given

- a set of packages of different weights, and
- A shipping pallet/container/box that has a maximum weight

Find

- combination of packages that has the maximum weight but not over the target

Solution:

- build all possible combinations by adding packages until over the target weight.
- remember the best combination

Combinations

• packet 0	yes	no
• packet 1	yes	no
• packet 2	yes	no
• packet 3	yes	no
• packet 4	yes	no
• packet 5	yes	no
• packet 6	yes	no
• packet 7	yes	no
• packet 8	yes	no
• packet 9	yes	no
• packet 10	yes	no
• packet 11	yes	no

Algorithm1

- Start with an empty combination
- initialise bestCombination and bestTotal to 0;
- Find combinations using additional packets from index 0.

- To find combinations using additional packets from index $i \dots$:
 - if including packet i would still be within target
 - add it to the current combination
 - if it beats the current best, then remember total and combination.
 - find combinations using additional packets from index $i+1 \dots$ < RECURSIVE CALL
 - remove it from the current combination
 - find combinations using additional packets from index $i+1 \dots$ < RECURSIVE CALL

Algorithm2

- use a binary integer to represent a combination:
10001110011010 => packets 1, 5, 6, 7, 10, 11,13
- Step through all numbers from 1 to 1111111111...111 to try all combinations
 - **for** combn from 1 to 11111111...1
 - work out total weight of combination
 - **if** weight <= target and weight > best so far
 - remember weight and combn

Cost of Algorithm2

- if n packets, then max combination represented by 2^n

- **for** combn from 1 to max

- work out total weight of combination
- **if** weight \leq target and weight $>$ best so far
 - remember weight and combn

with n packets, max = 2^n

$O(n)$

$O(1)$

$O(1)$



2^n times

Cost of Algorithm1

- Start with an empty combination
- initialise bestCombination and bestTotal to 0;
- Find combinations (target, packets, combn, 0, 0).

- To find combinations (target, packets, combSoFar, index, totalSoFar)
 - packet = packets[index]
 - newTotal= totalSoFar+ packet.weight
 - if newTotal <= target
 - add packet to the combSoFar
 - **if** newTotal > bestTotal then bestTotal=newTotal, bestCombination = combSoFar
 - find combinations(target, packets combSoFar, index+1, newTotal)
 - remove packet from combSoFar
 - find combinations(target, packets, combSoFar, index+1, totalSoFar)

Cost of Algorithm 1

- $\text{Cost}(n)$ = cost of finding with n remaining packets to try
- $\text{Cost}(1) = O(1)$
- $\text{Cost}(n) = O(1) + \text{Cost}(n-1) + \text{Cost}(n-1)$
 $= 2 \text{Cost}(n-1) + O(1)$
 $= 2(2\text{Cost}(n-2) + O(1)) + O(1)$

The cost approximately doubles when n increase by 1 $\Rightarrow O(2^n)$